

# Logical relations for fine-grained concurrency

Aaron Turon

*with*

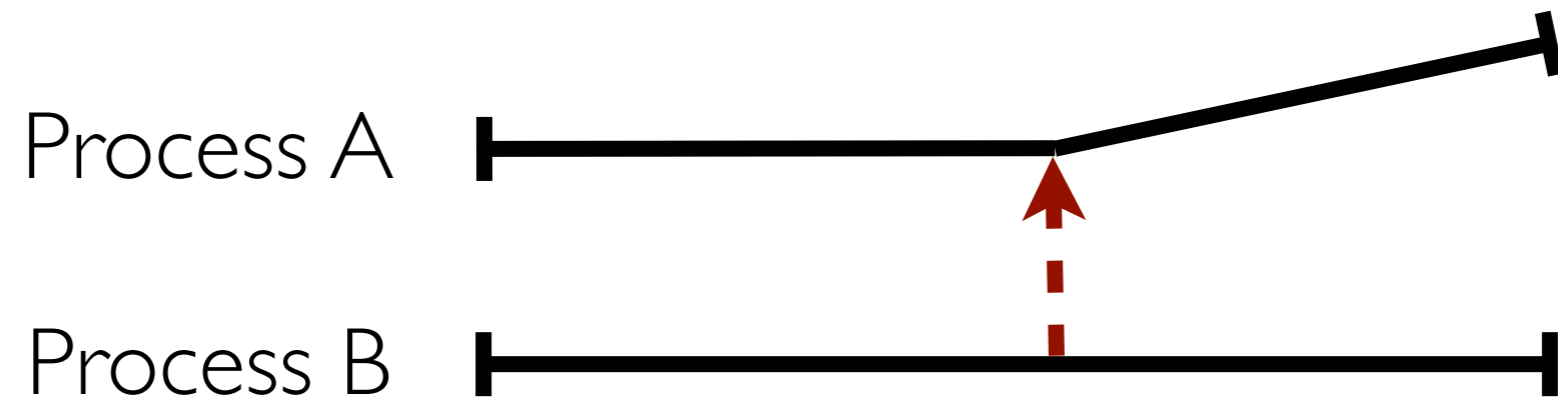
Jacob Thamsborg, Amal Ahmed,  
Lars Birkedal, Derek Dreyer

**Concurrency** is overlapped execution:

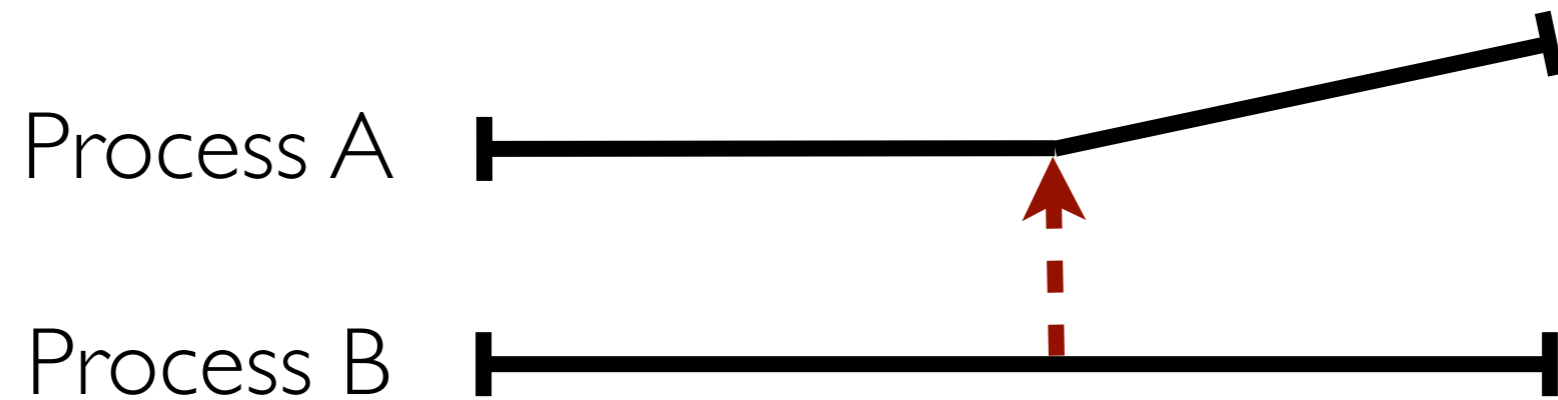
Process A 

Process B 

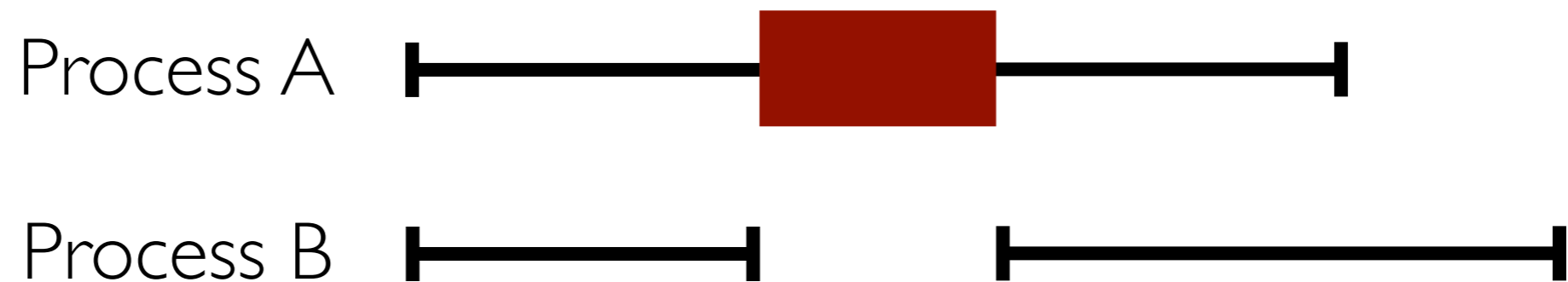
**Concurrency** is overlapped execution:



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**Atomicity** is its antidote:



# Granularity (in time & space)



**Coarse-grained** atomicity:  
easier reasoning



**Fine-grained** atomicity:  
lower latency  
higher throughput (if parallel)

# Atomicity abstraction



Fine-grained  
implementation



Coarse-grained  
specification

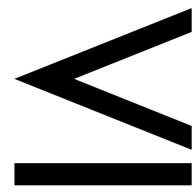
# Atomicity abstraction




Fine-grained  
implementation



Coarse-grained  
specification



  
Refines, i.e.,  
contextually  
approximates

# Contributions

*Direct, insightful* refinement method  
(no linearizability!)

Handles *higher-order, polymorphic* languages

Scales to *sophisticated* algorithms



# Key Idea

**Atomicity abstraction**, like data abstraction, relies on hiding

# Approach

Transition systems  
for hidden state

} *State of the art:  
Kripke Logical  
Relations, ICFP 10*

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Transition systems  
for hidden state

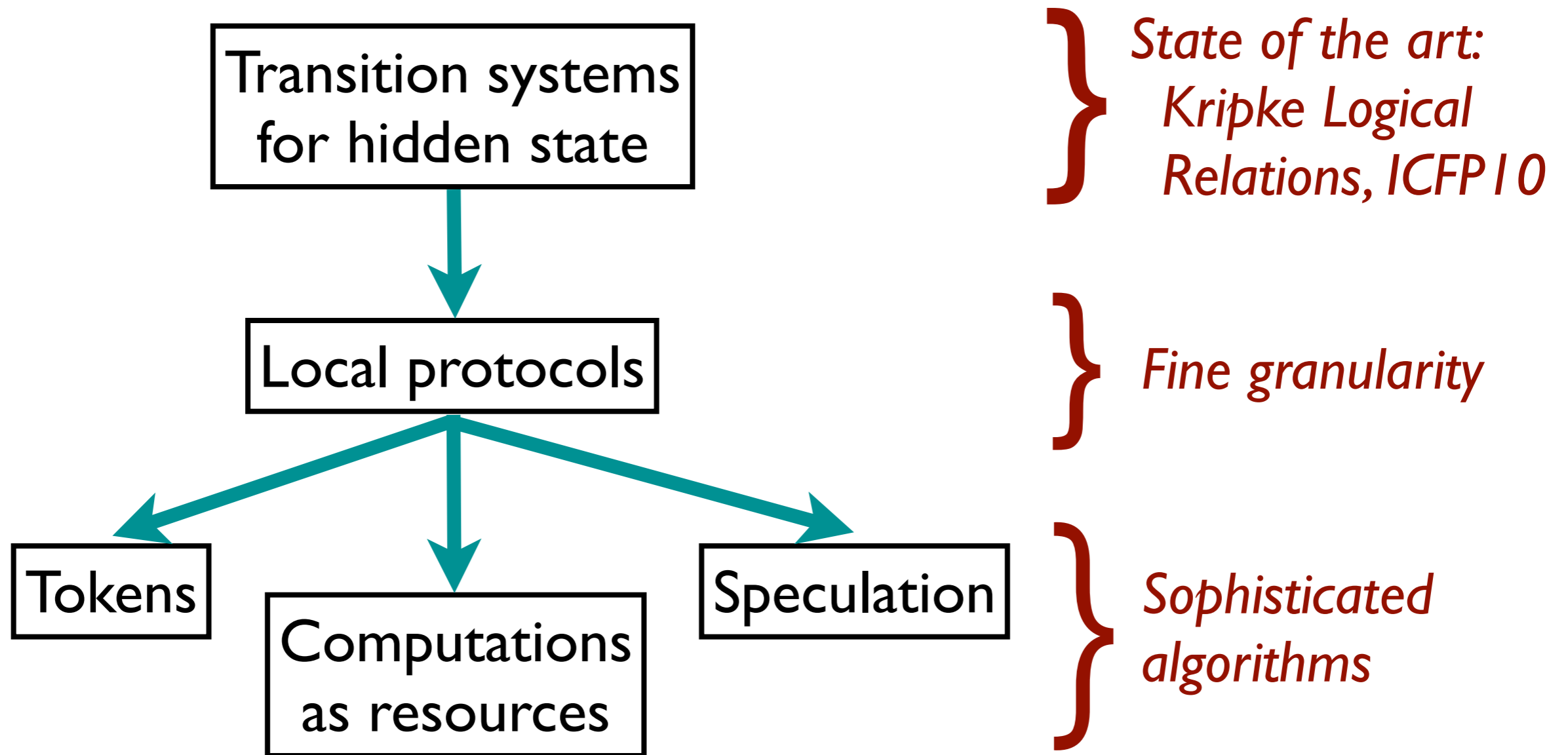


Local protocols

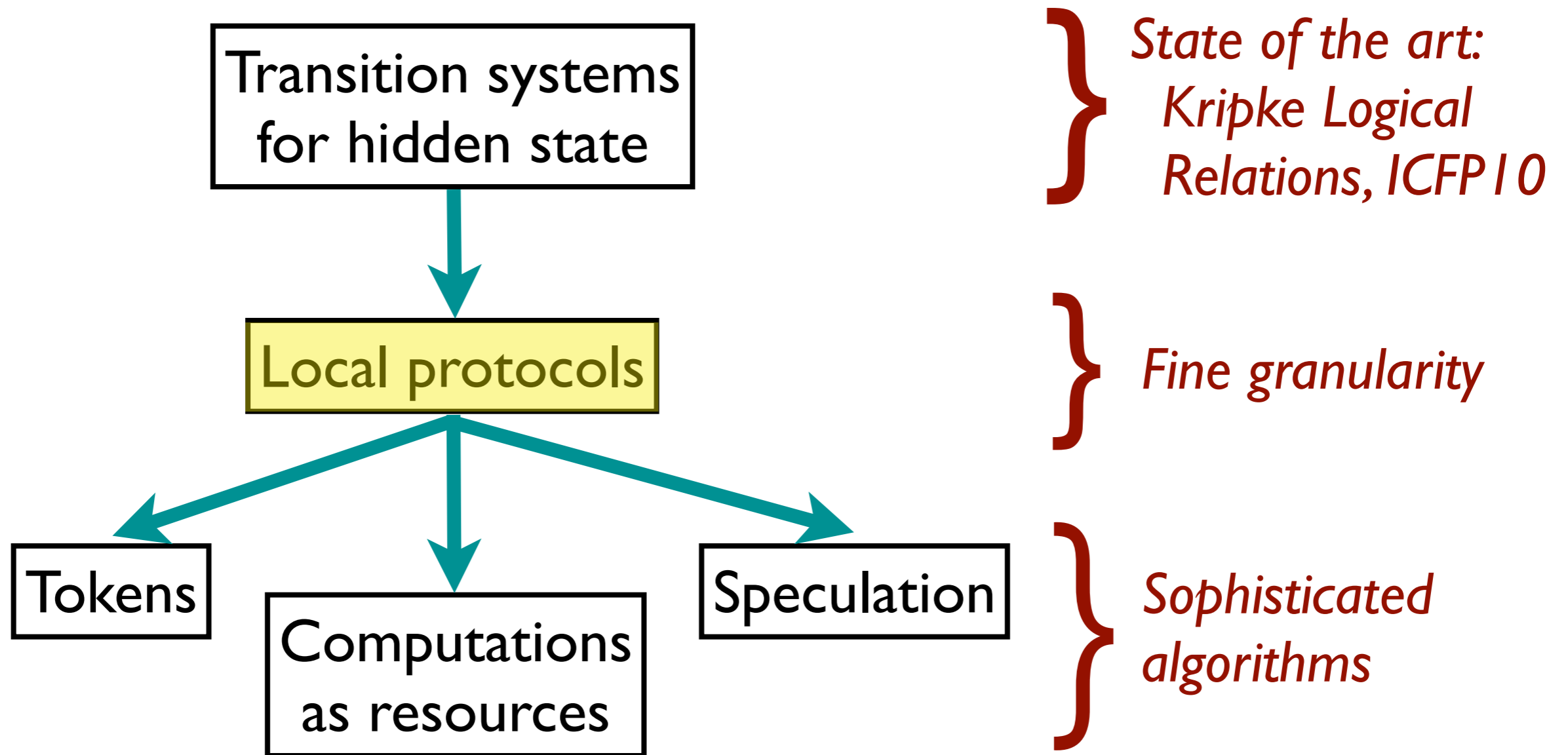
} *State of the art:  
Kripke Logical  
Relations, ICFP 10*

} *Fine granularity*

# Approach



# Approach





# Compare and set

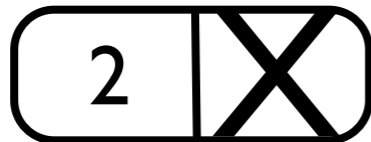
```
cas(cell, old, new) = atomic {  
  if !cell := old  
  then cell := new; true  
  else false  
}
```

# A Lock-free “Queue”

Head 

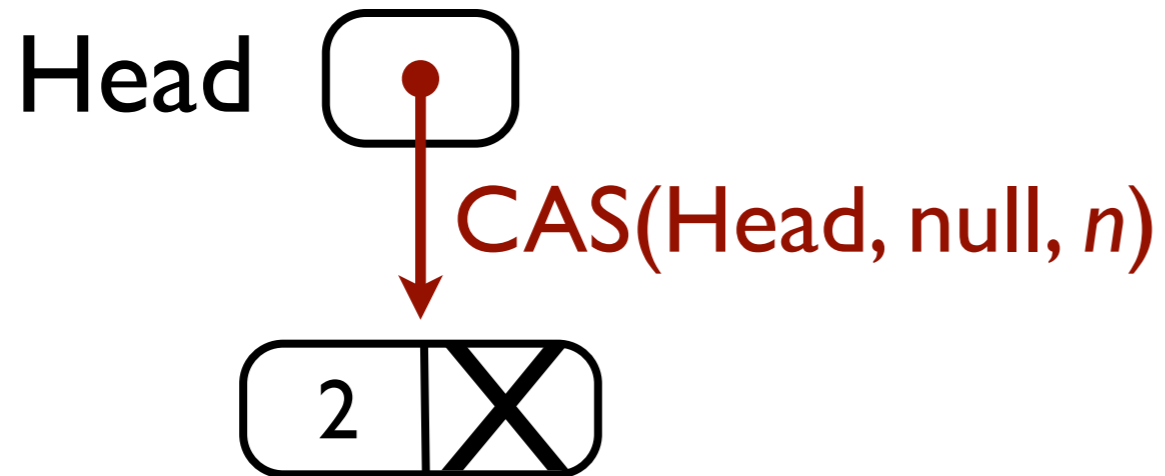
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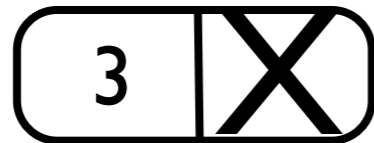
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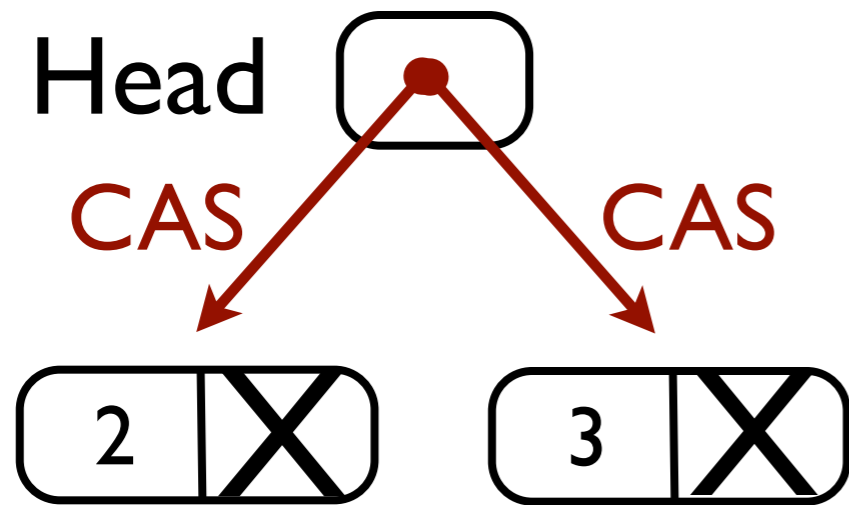
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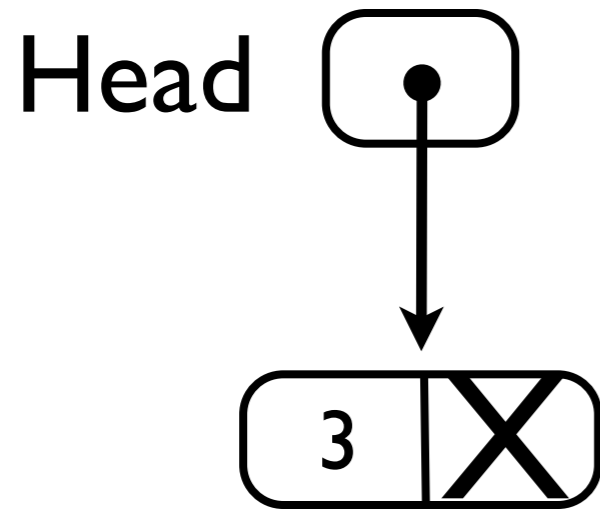




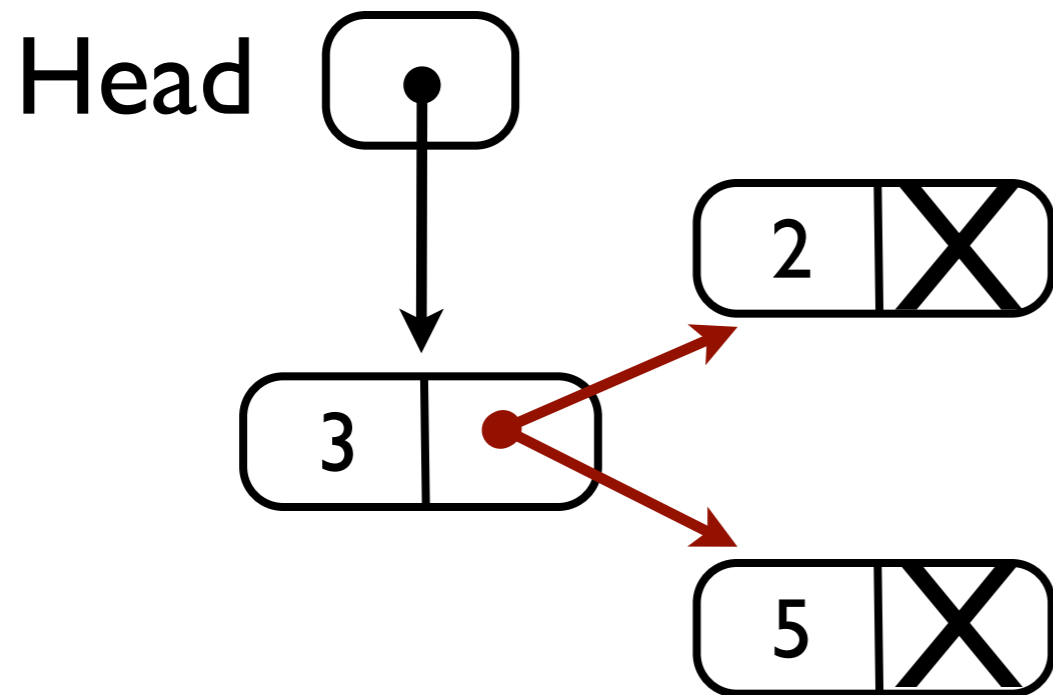
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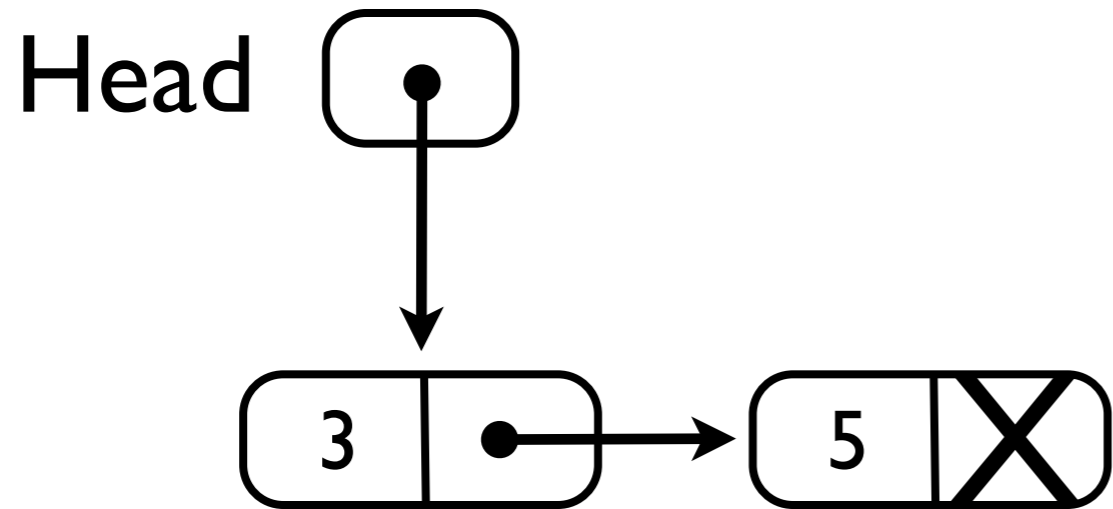
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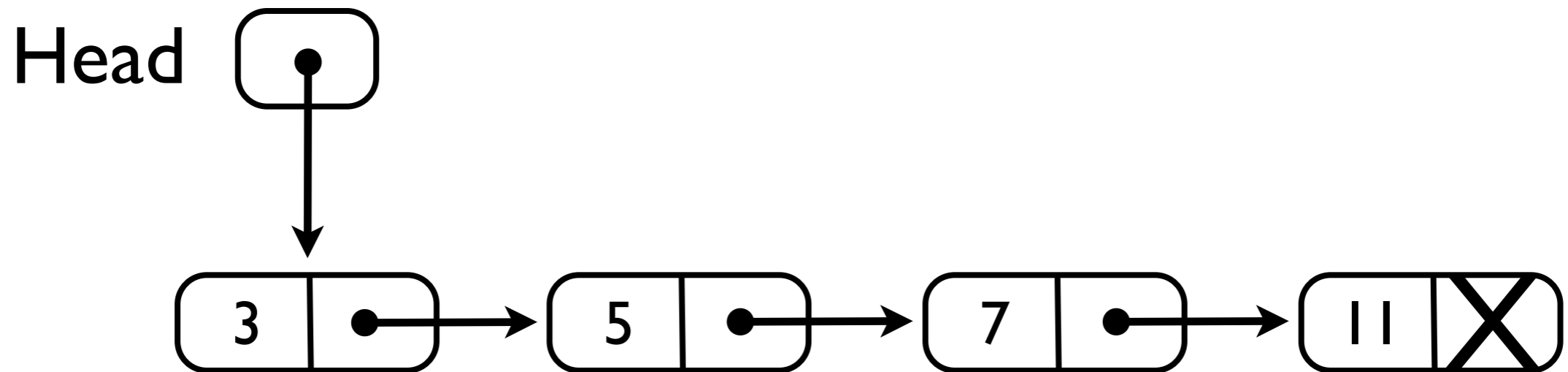
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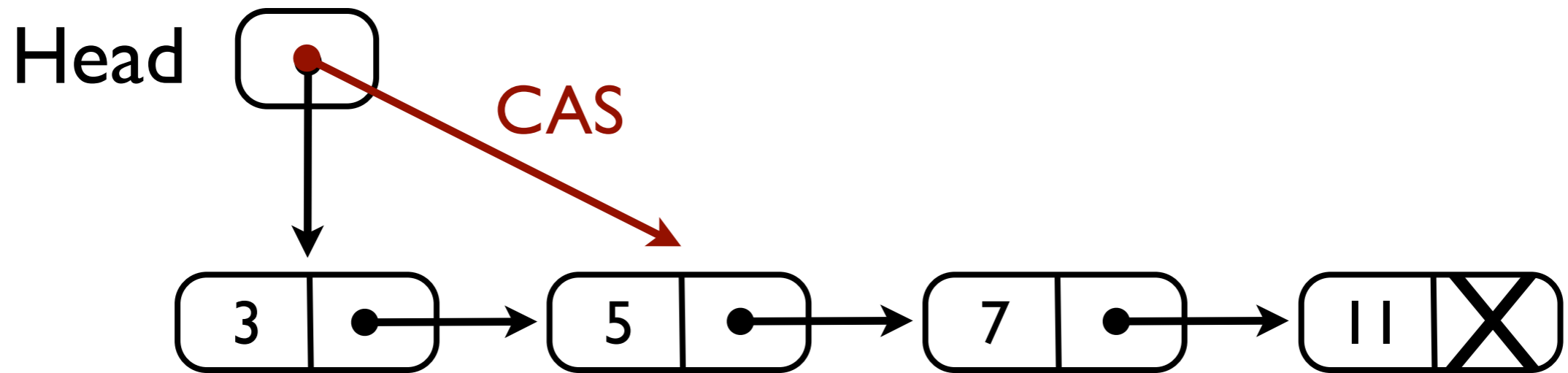


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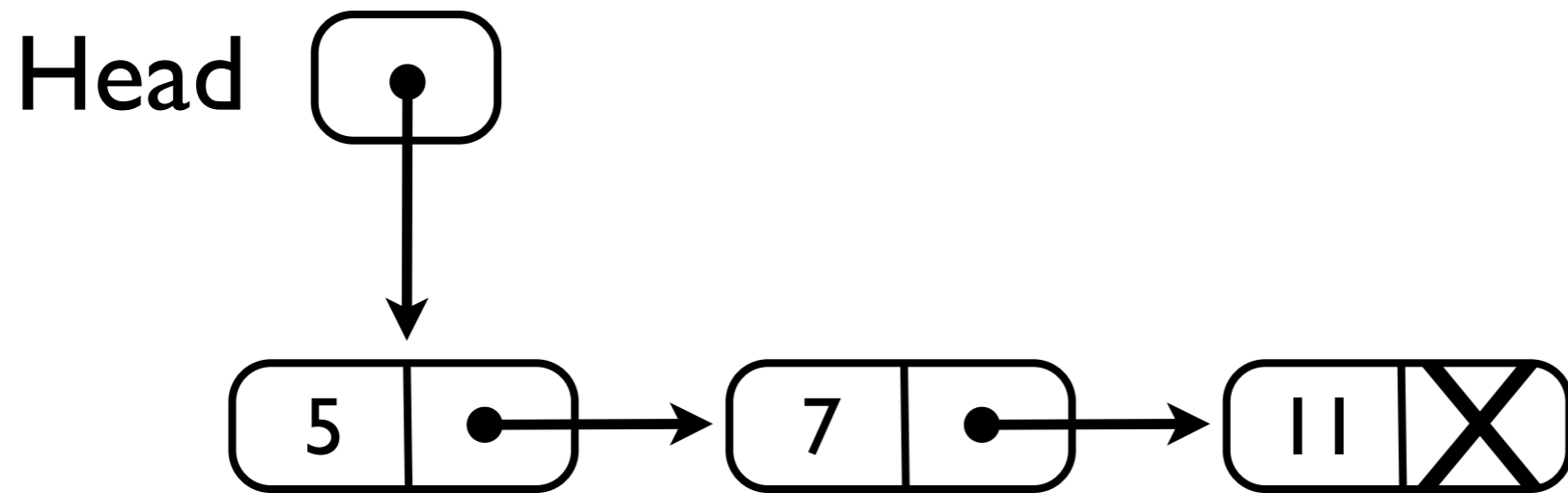




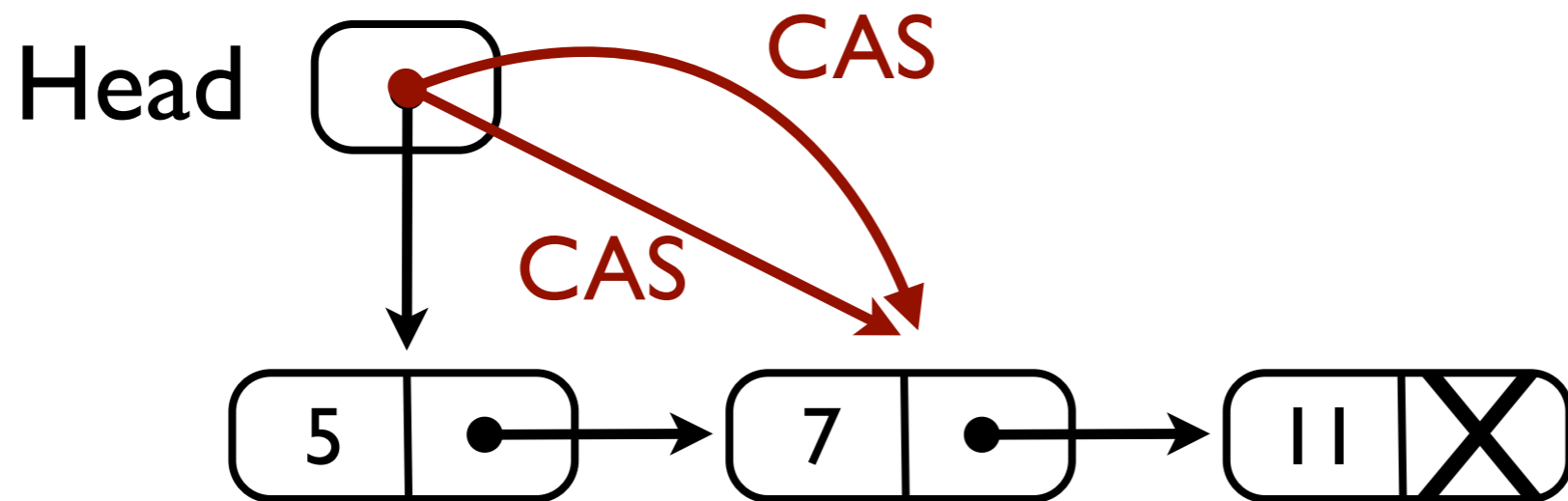
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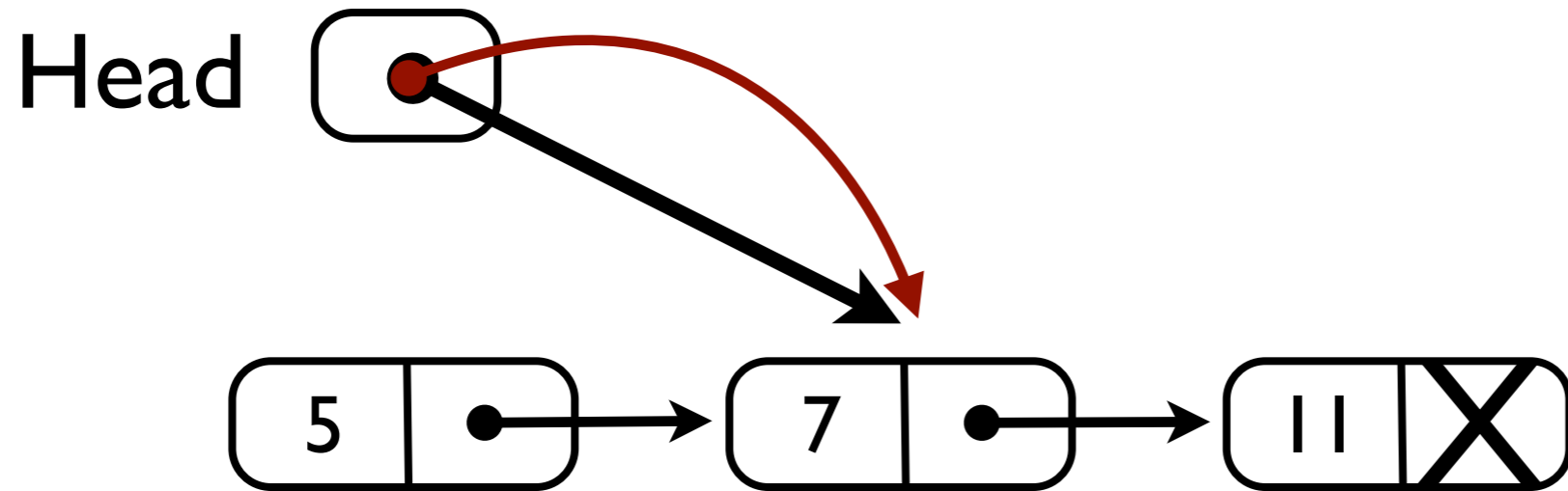
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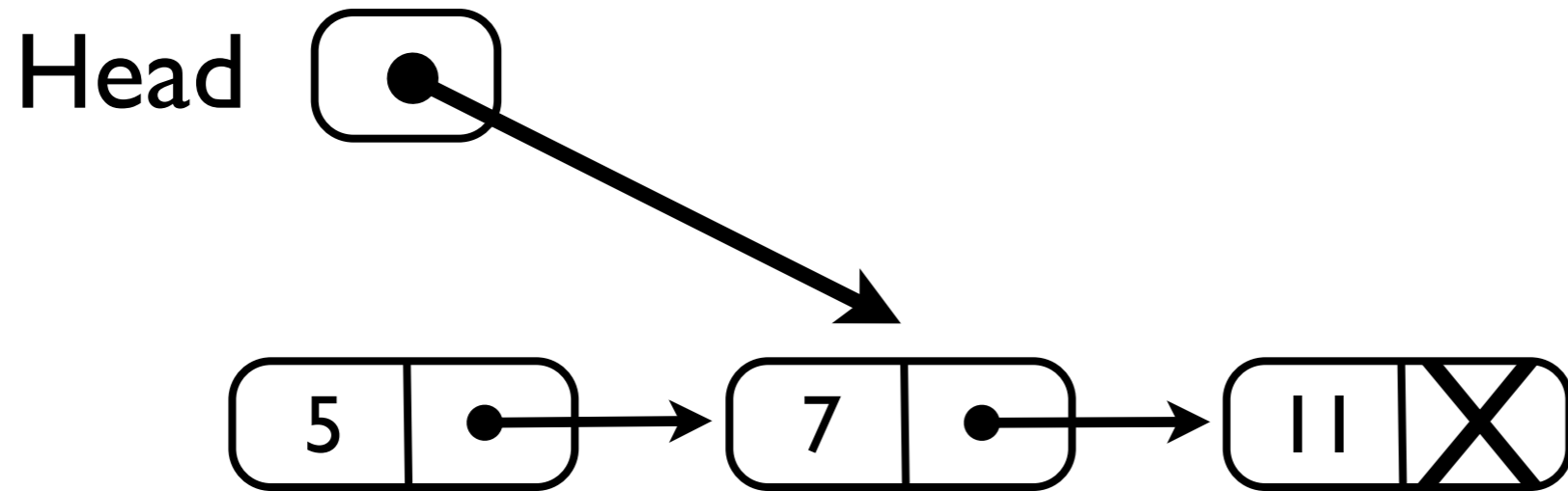
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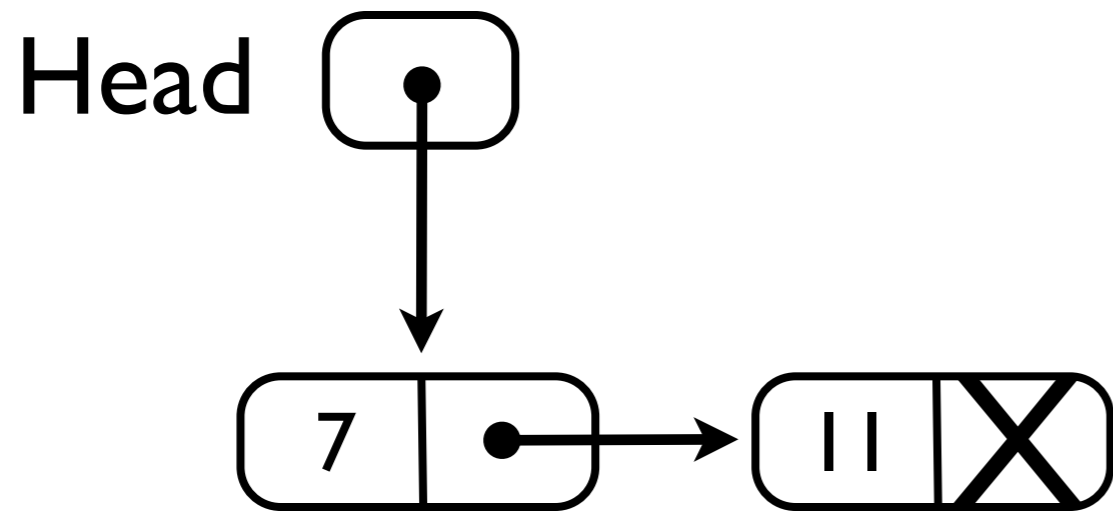
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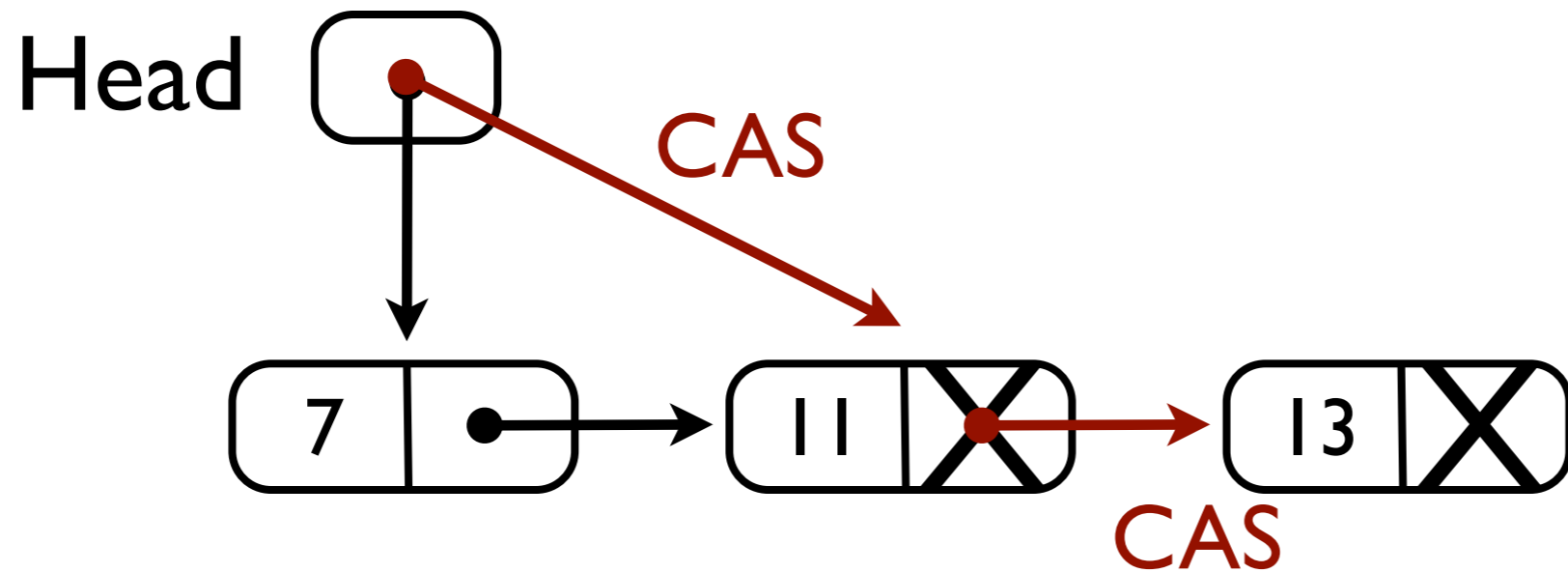
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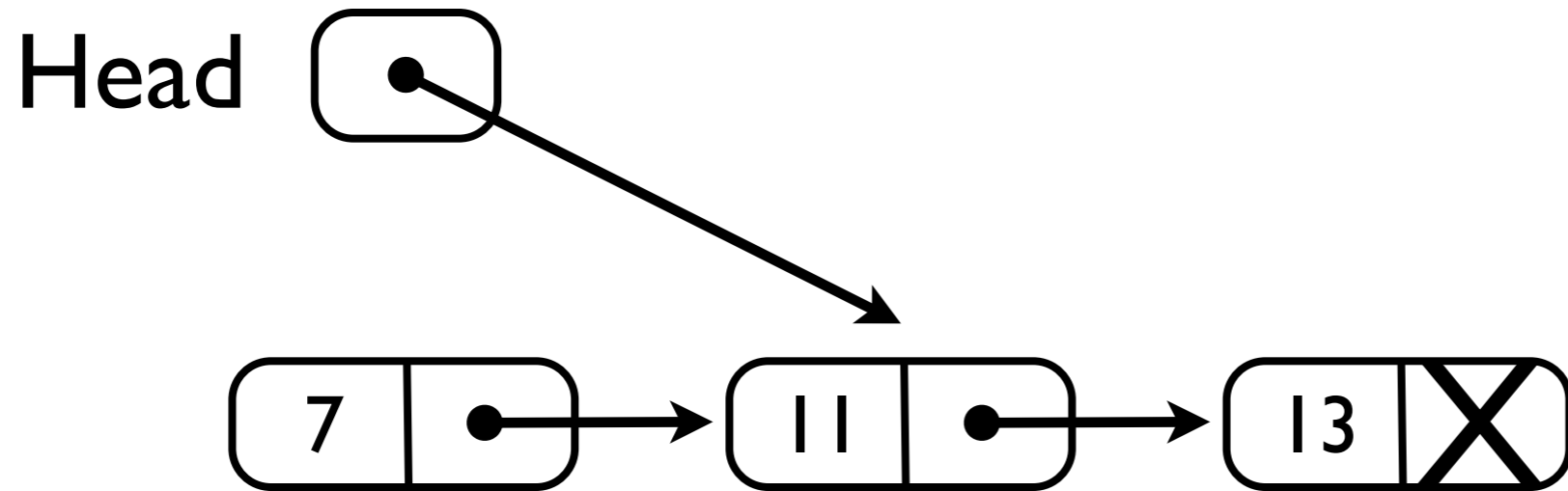
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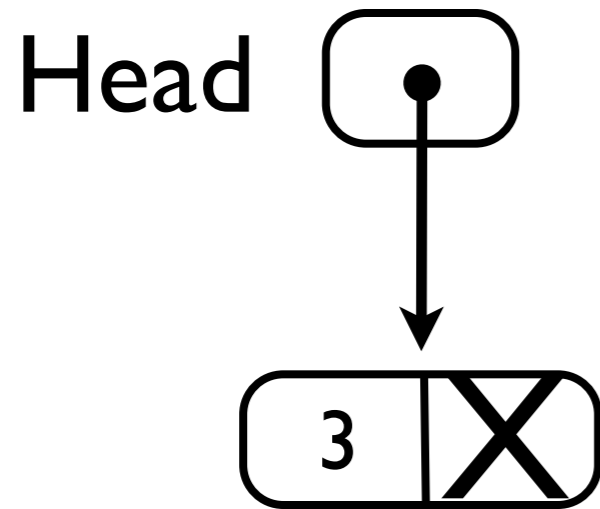


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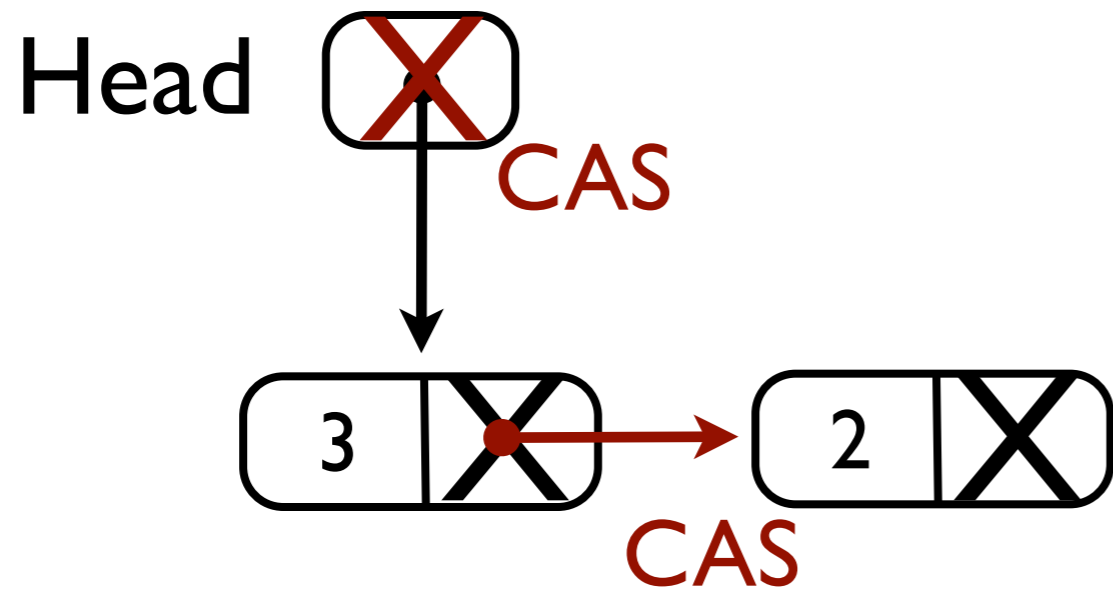




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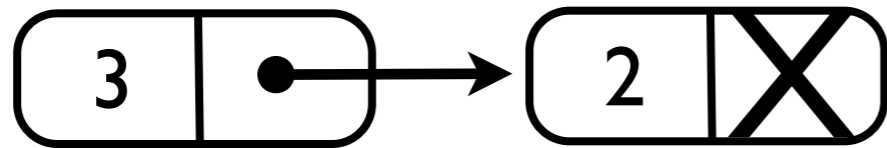


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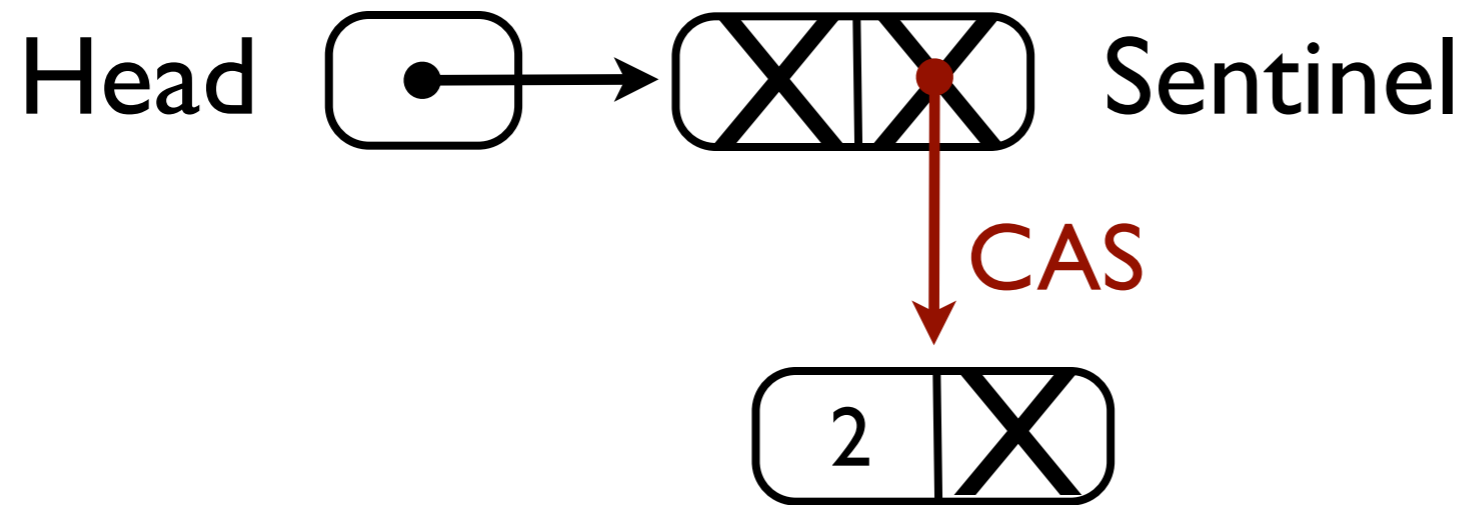
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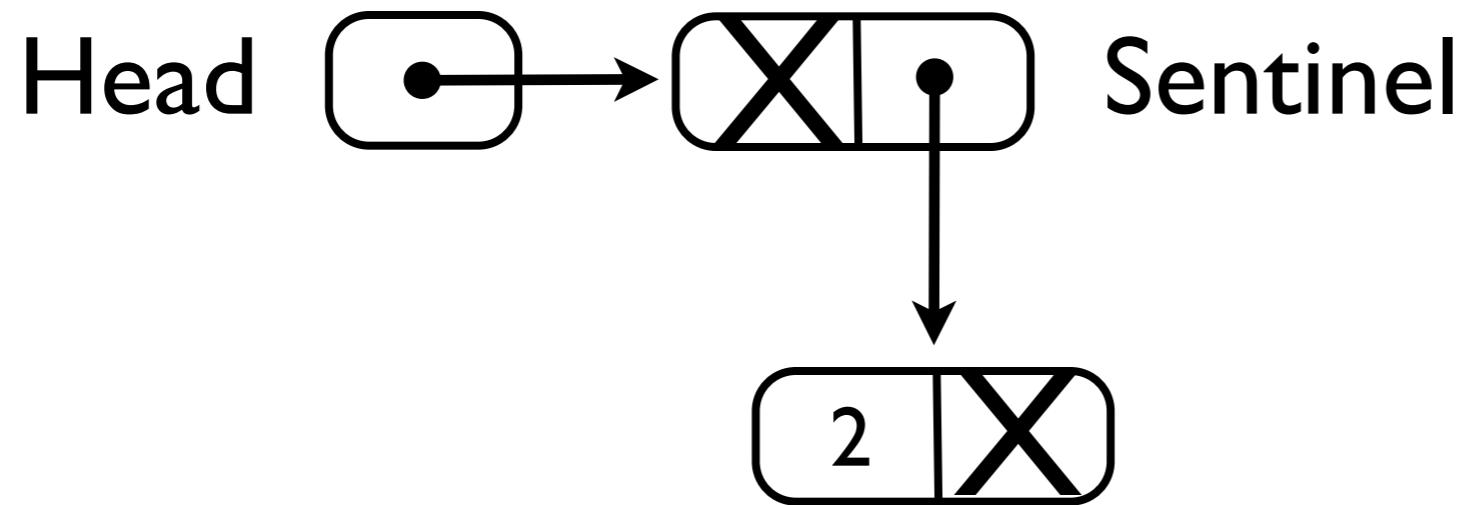
# The Michael-Scott Queue



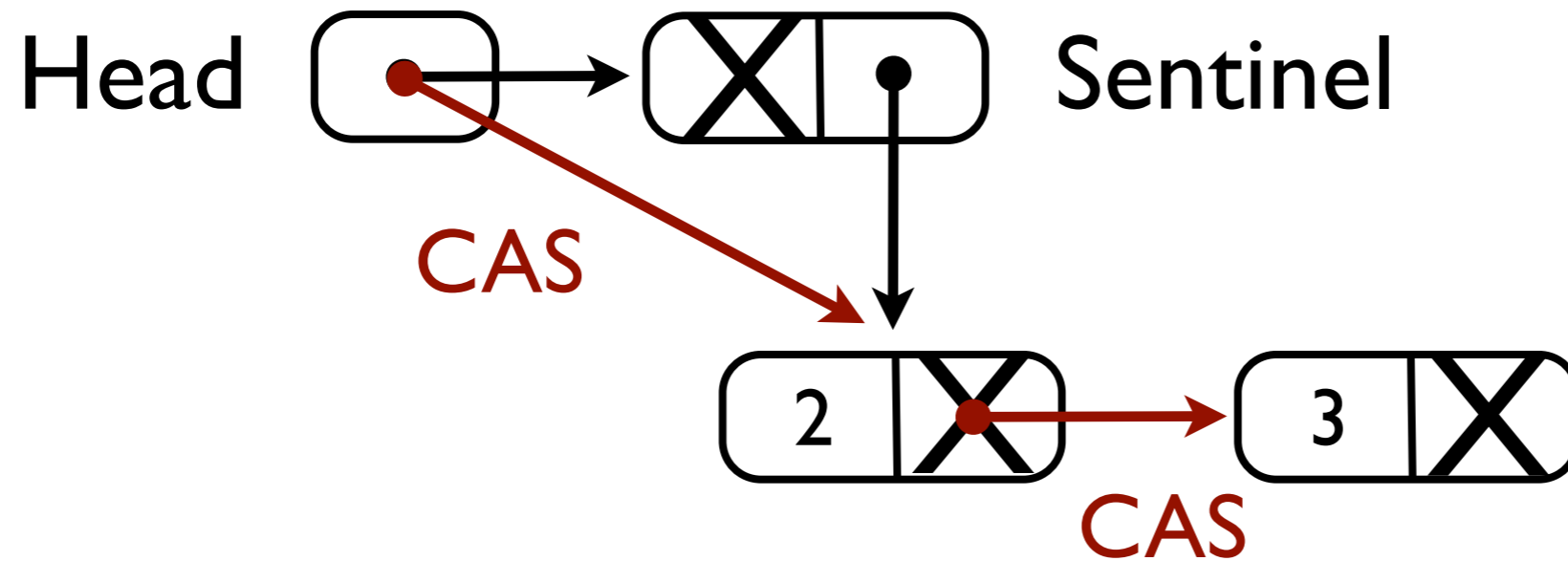
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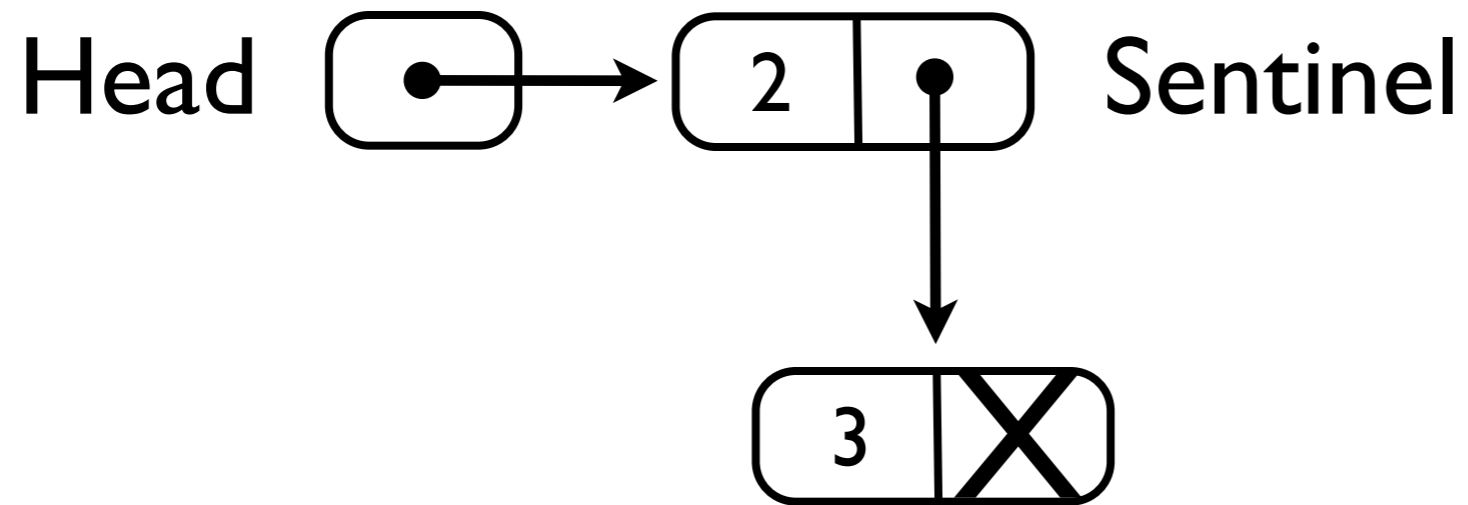
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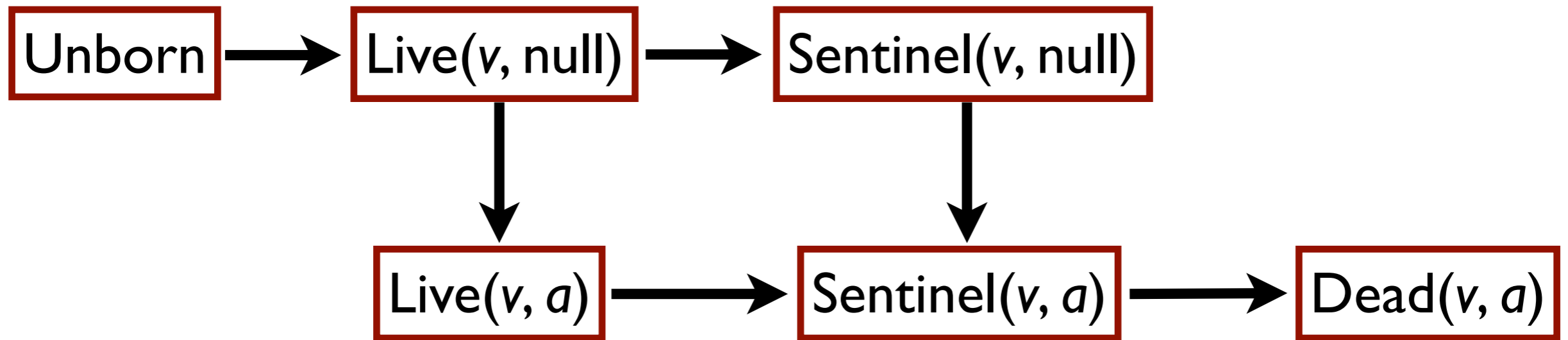


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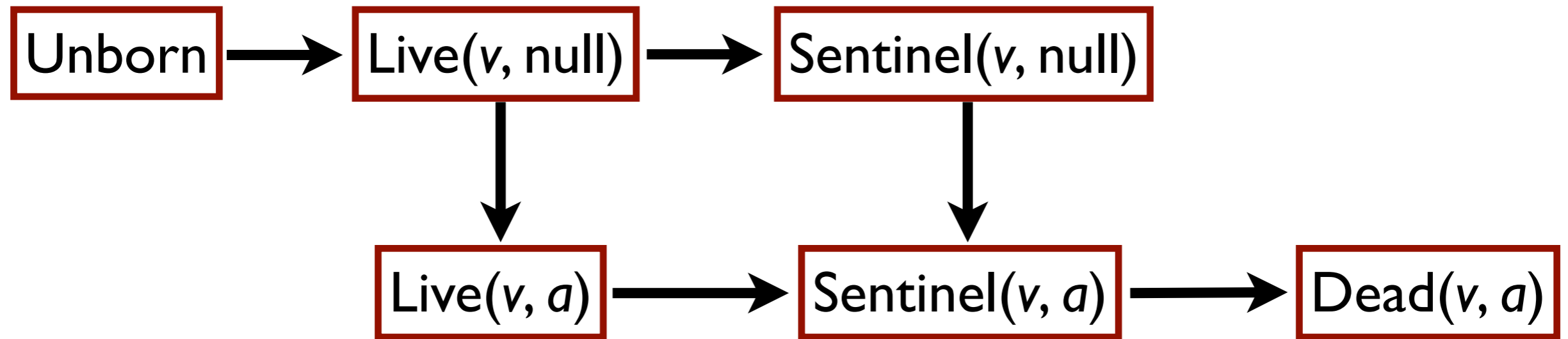




# The life story of a node



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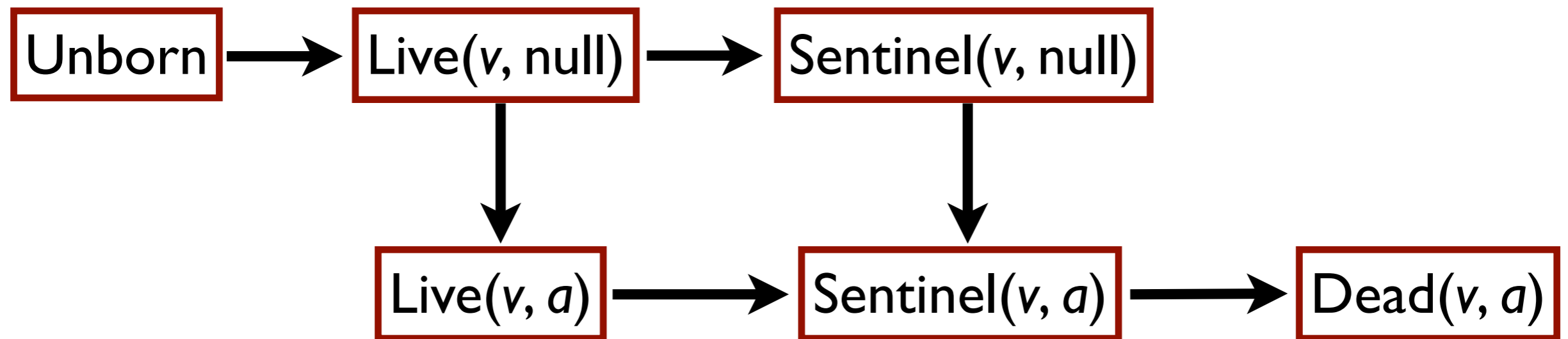
## **Global constraints:**

Exactly one sentinel

Live iff reachable from sentinel

Nonnull next pointers are born

# The life story of a node



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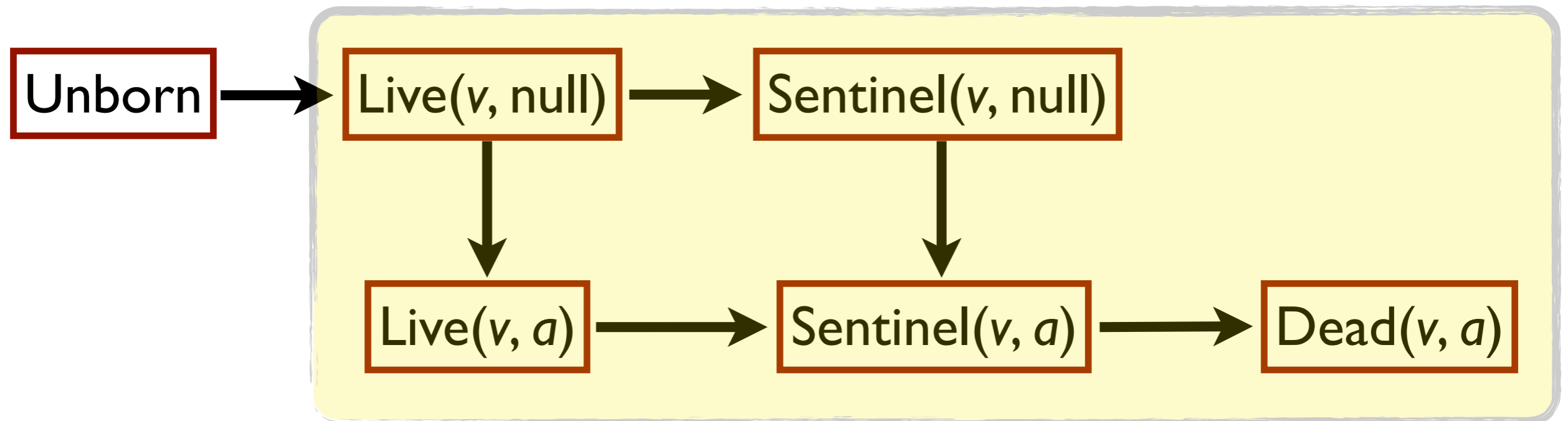
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linkNode(p, n) = case !(p.next)
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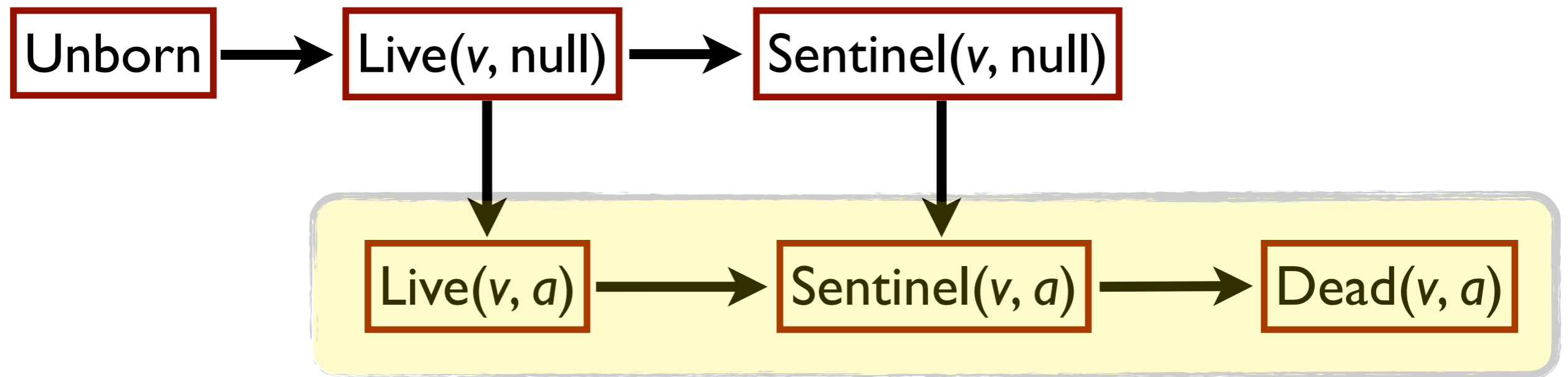
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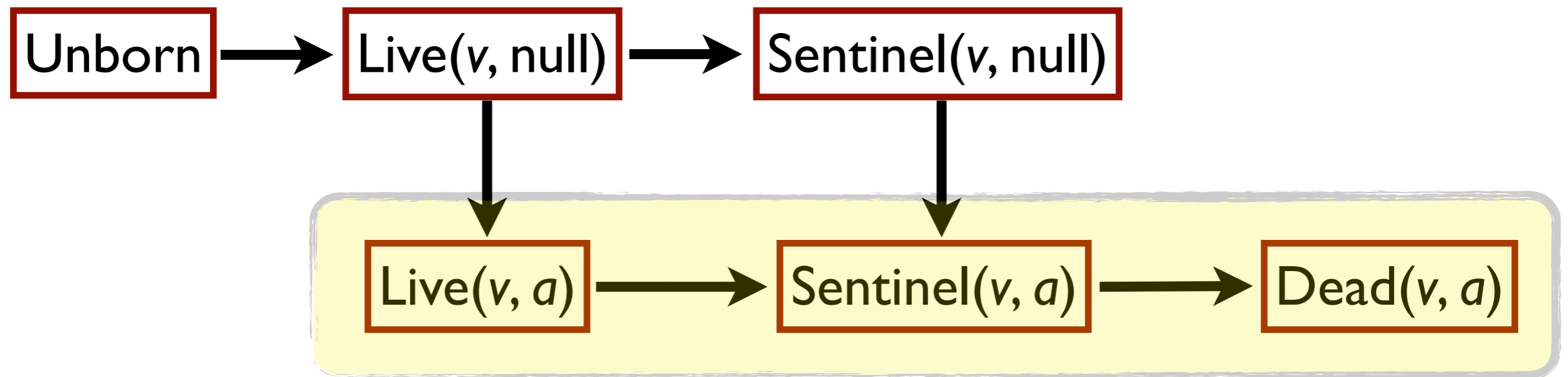
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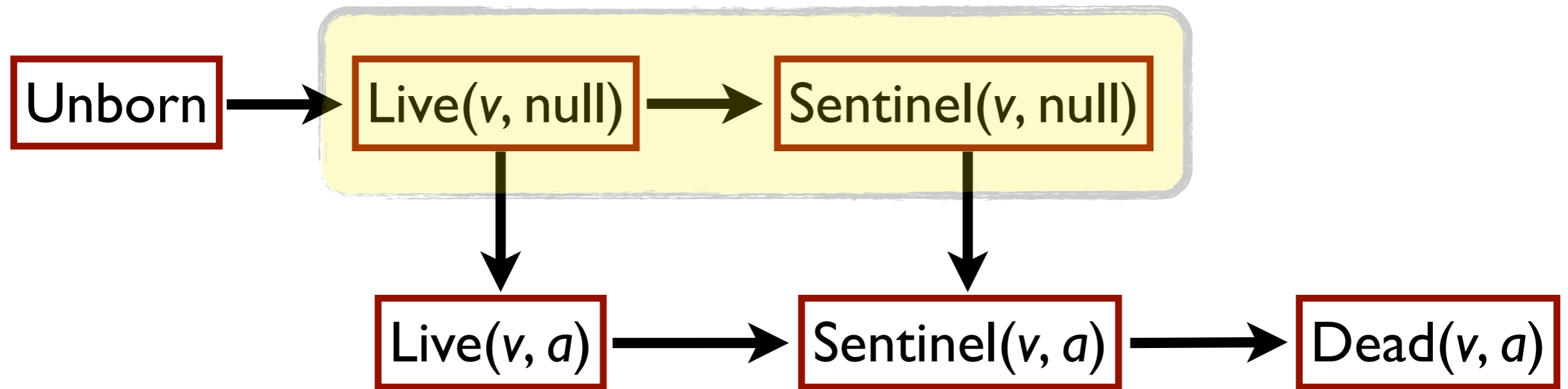
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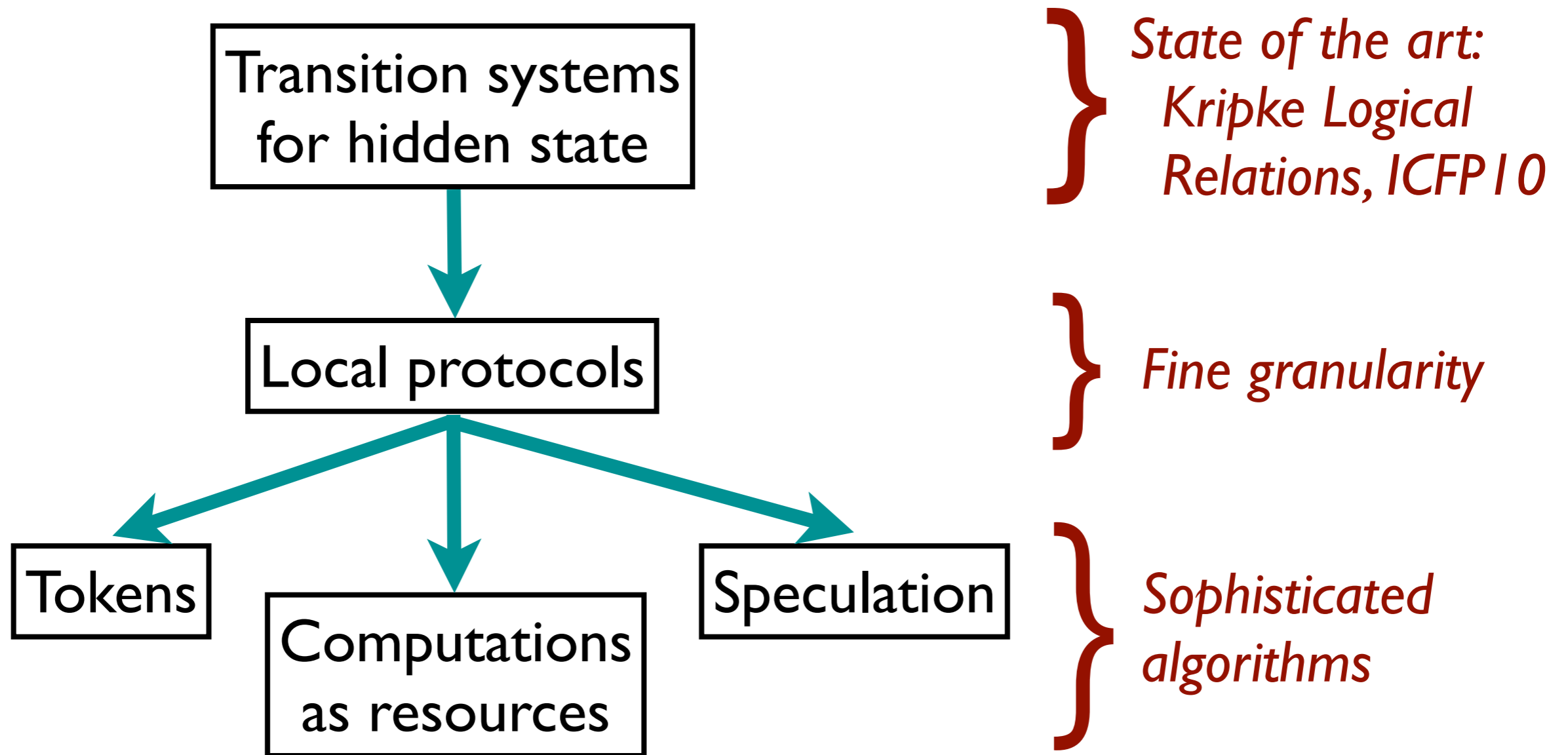
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# Key Idea

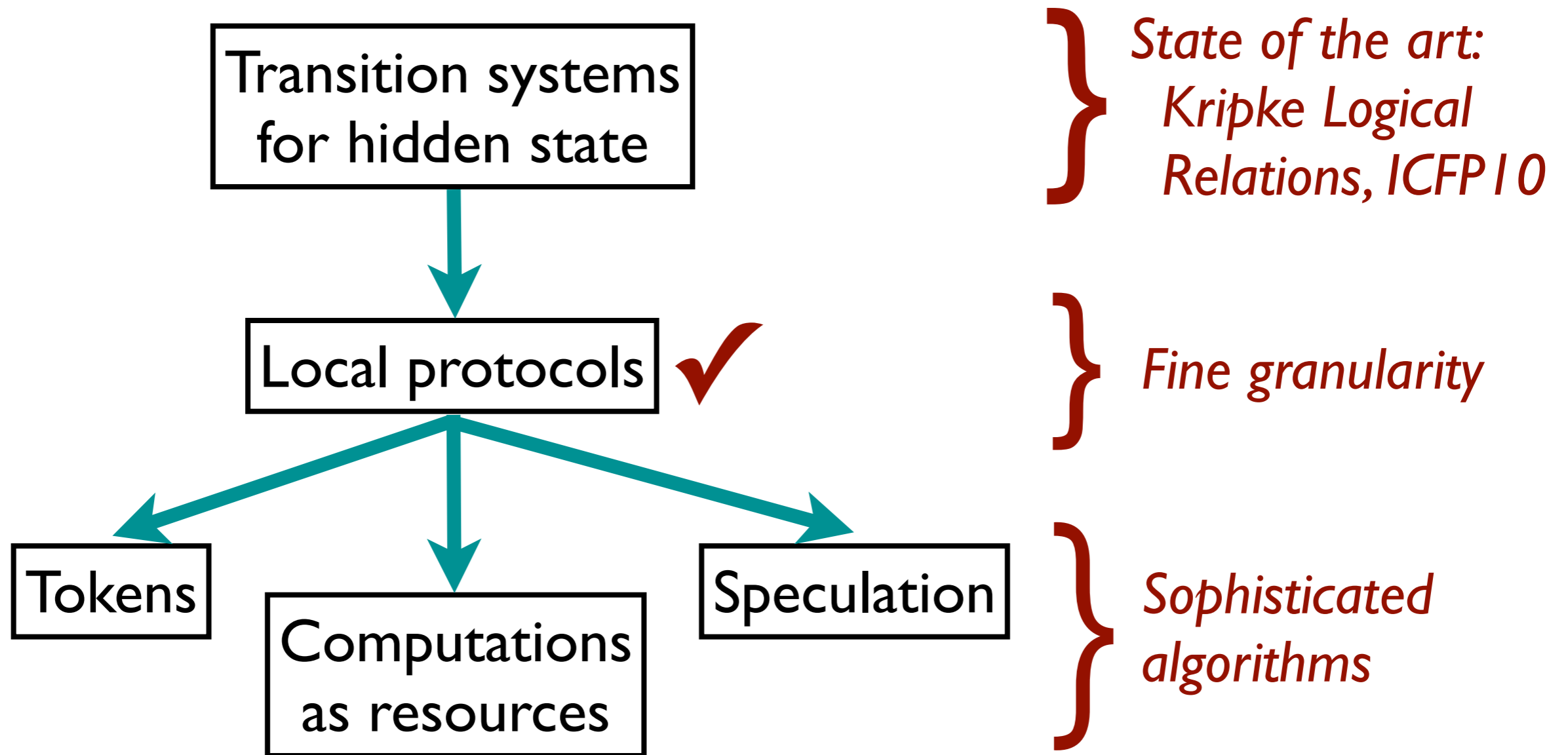
**Local protocols** capture abstract knowledge/interference at the same granularity as the algorithm operates



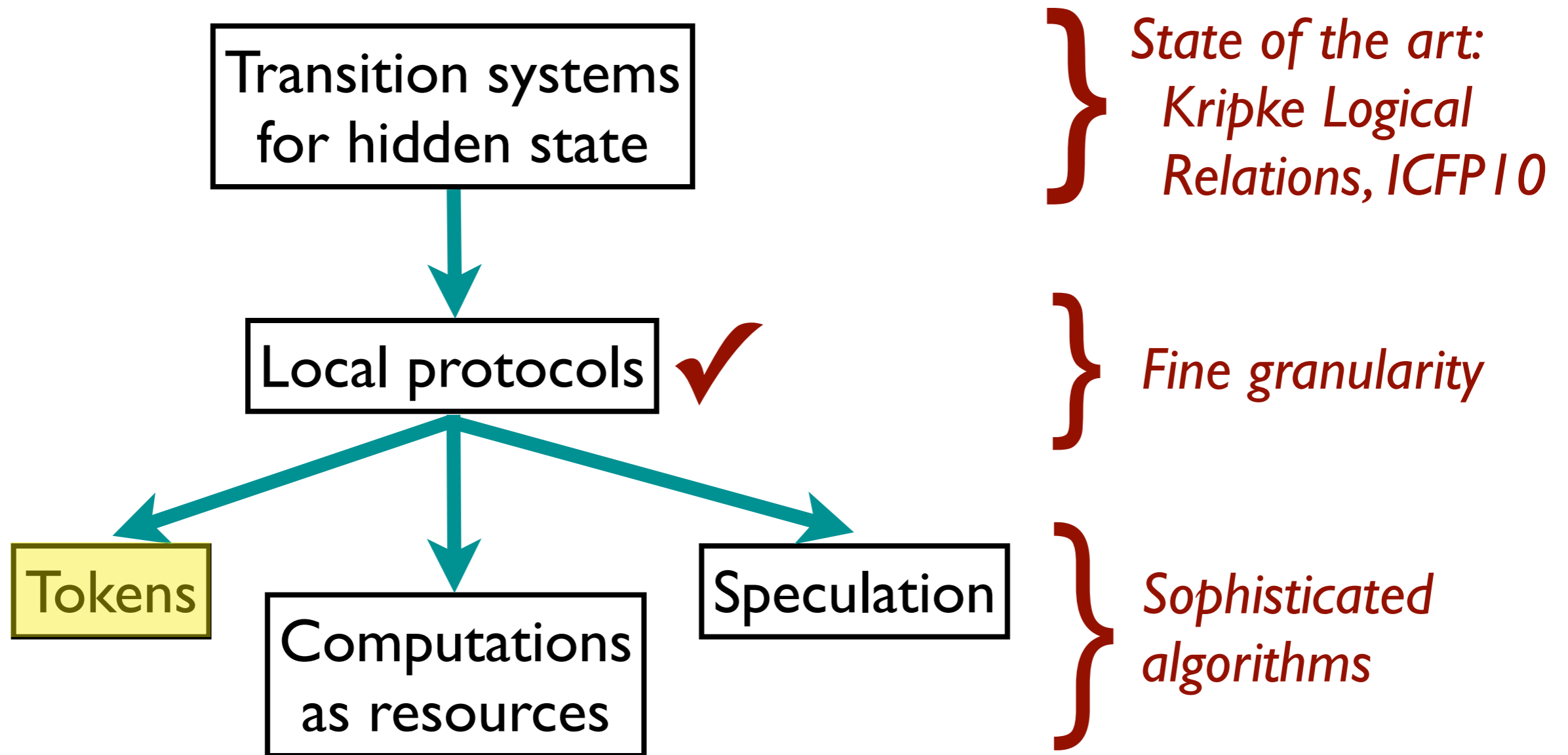
# Approach



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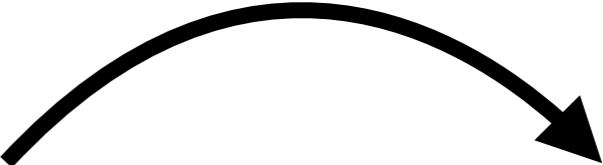


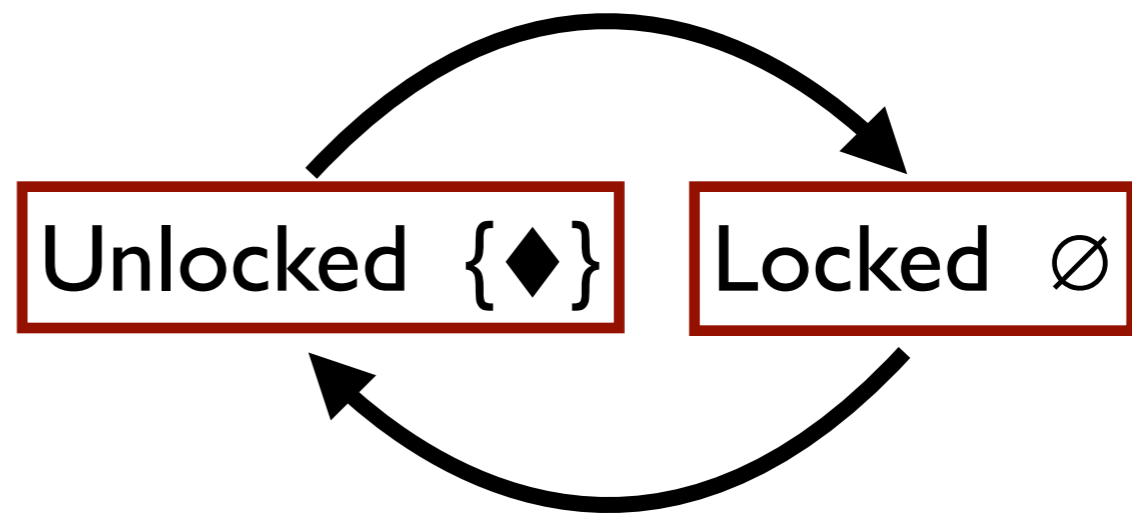
# Approach



Unlocked

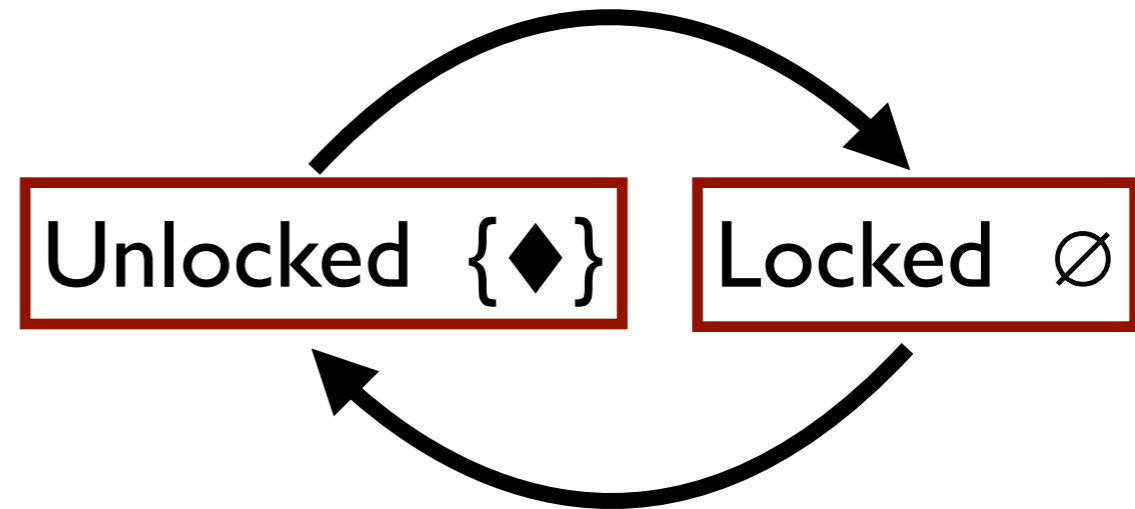
Locked





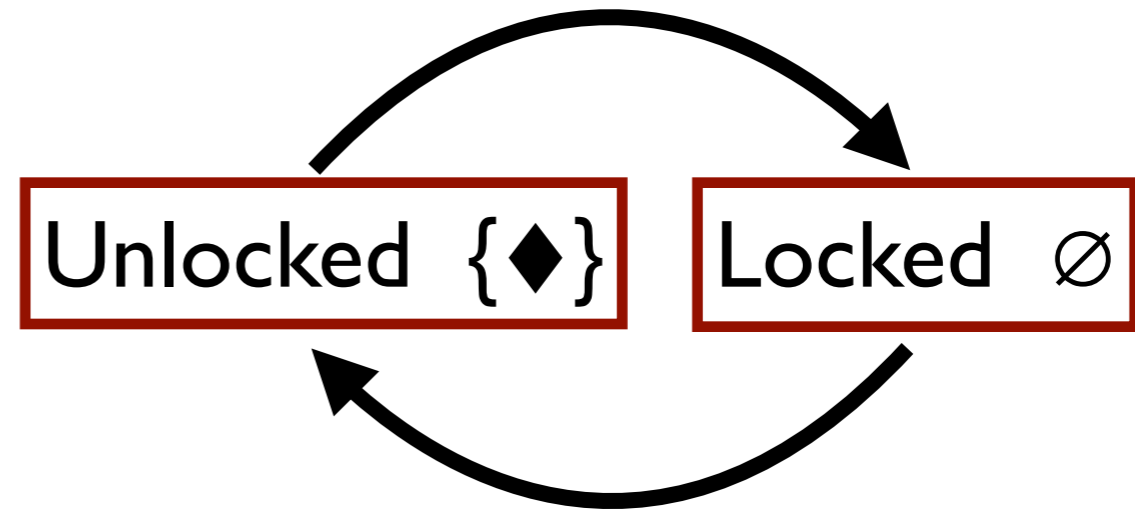
# Conservation Law

$$\begin{array}{ccc} \textit{Before} & & \textit{After} \\ T \cup P & = & T' \cup P' \\ \uparrow & & \uparrow \\ \text{Thread's} & & \text{Protocol's} \\ \text{tokens} & & \text{tokens} \end{array}$$



# Conservation Law

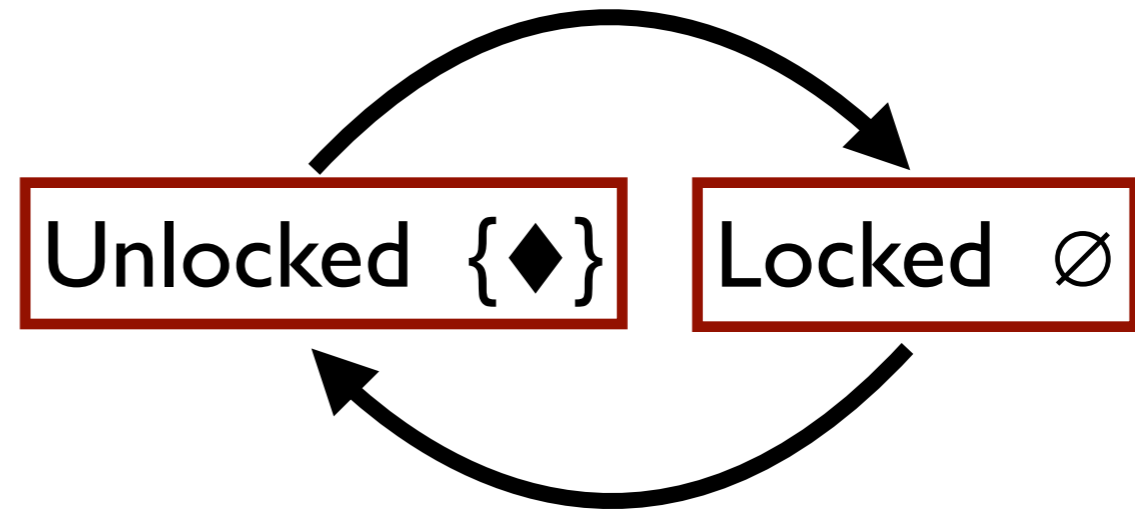
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Unlocked → Locked:  $\emptyset \cup \{\diamond\} = ? \cup \emptyset$

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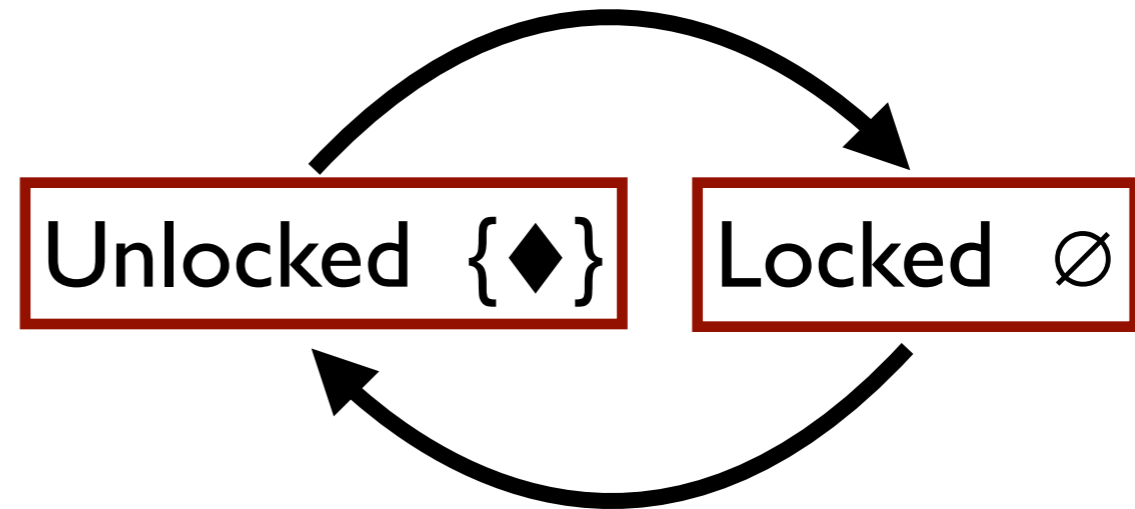


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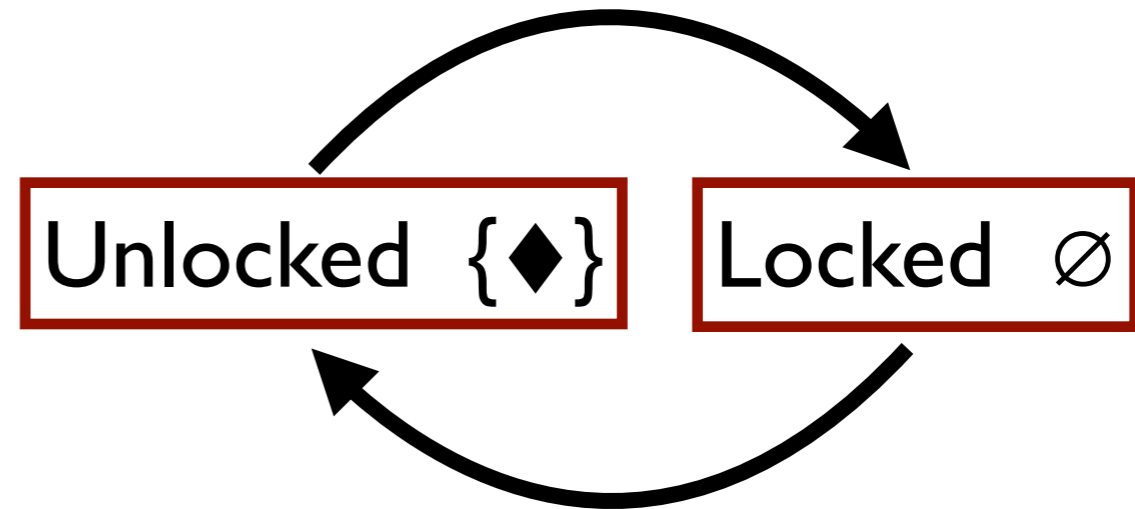
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$$\text{Unlocked} \rightarrow \text{Locked: } \begin{array}{cccc} T & P & T' & P' \\ \emptyset & \uplus \{ \blacklozenge \} & = & \{ \blacklozenge \} \uplus \emptyset \end{array}$$

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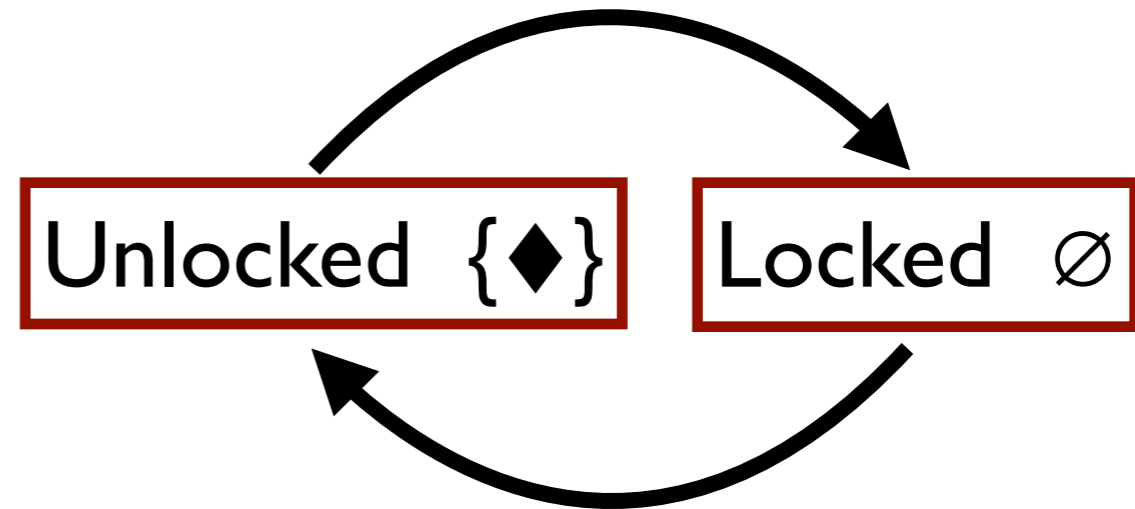


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$$\text{Locked} \rightarrow \text{Unlocked: } \begin{array}{cccc} T & P & T' & P' \\ \emptyset & \cup \emptyset & = & \text{?} \cup \{ \blacklozenge \} \end{array}$$

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$$\text{Unlocked} \rightarrow \text{Locked: } \emptyset \cup \{\diamond\} = \{\diamond\} \cup \emptyset$$

~~Locked  $\rightarrow$  Unlocked:  $\emptyset \cup \emptyset = ? \cup \{\diamond\}$  *Impossible*~~

# Key Idea

**Tokens** enable threads to gain and lose roles *dynamically*, via transitions

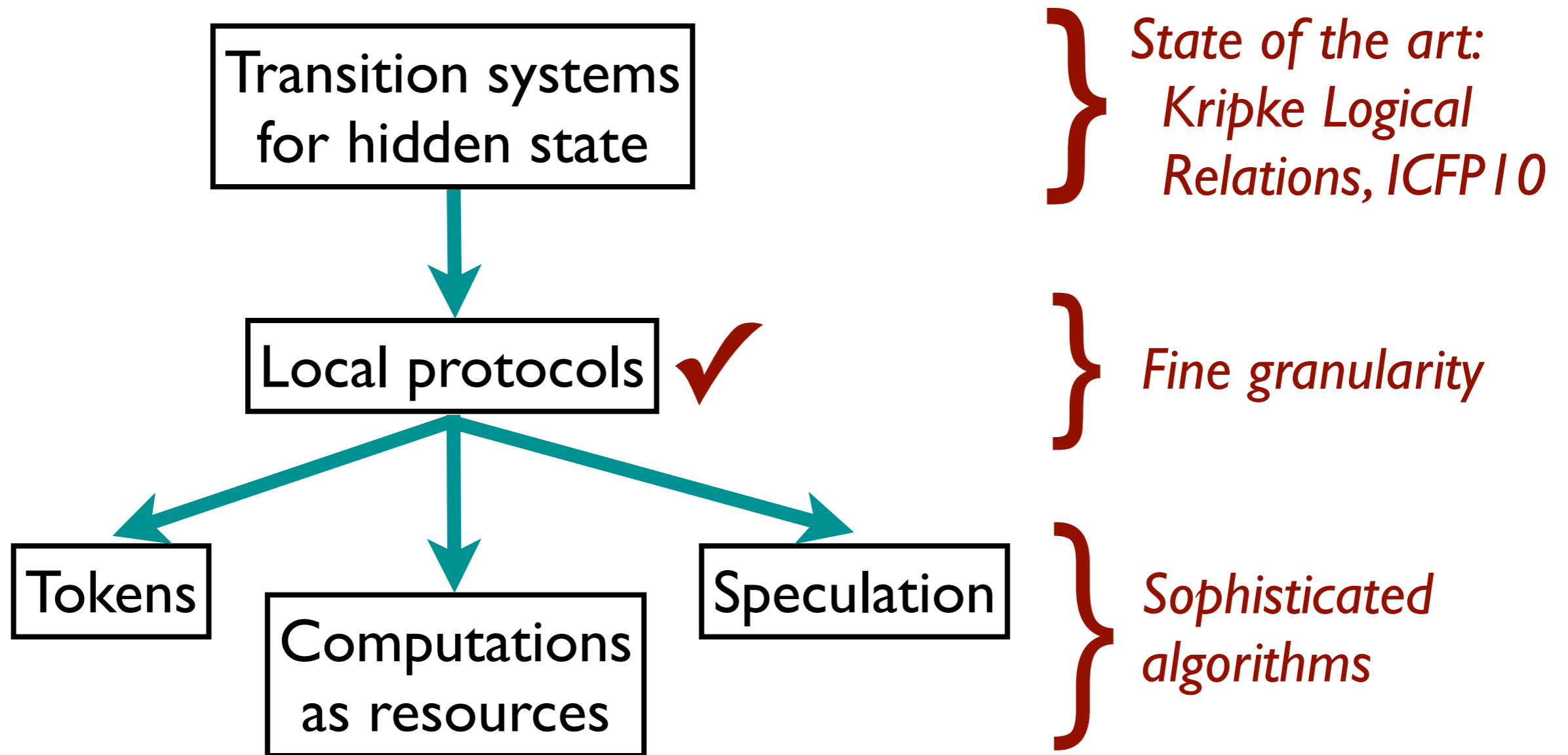
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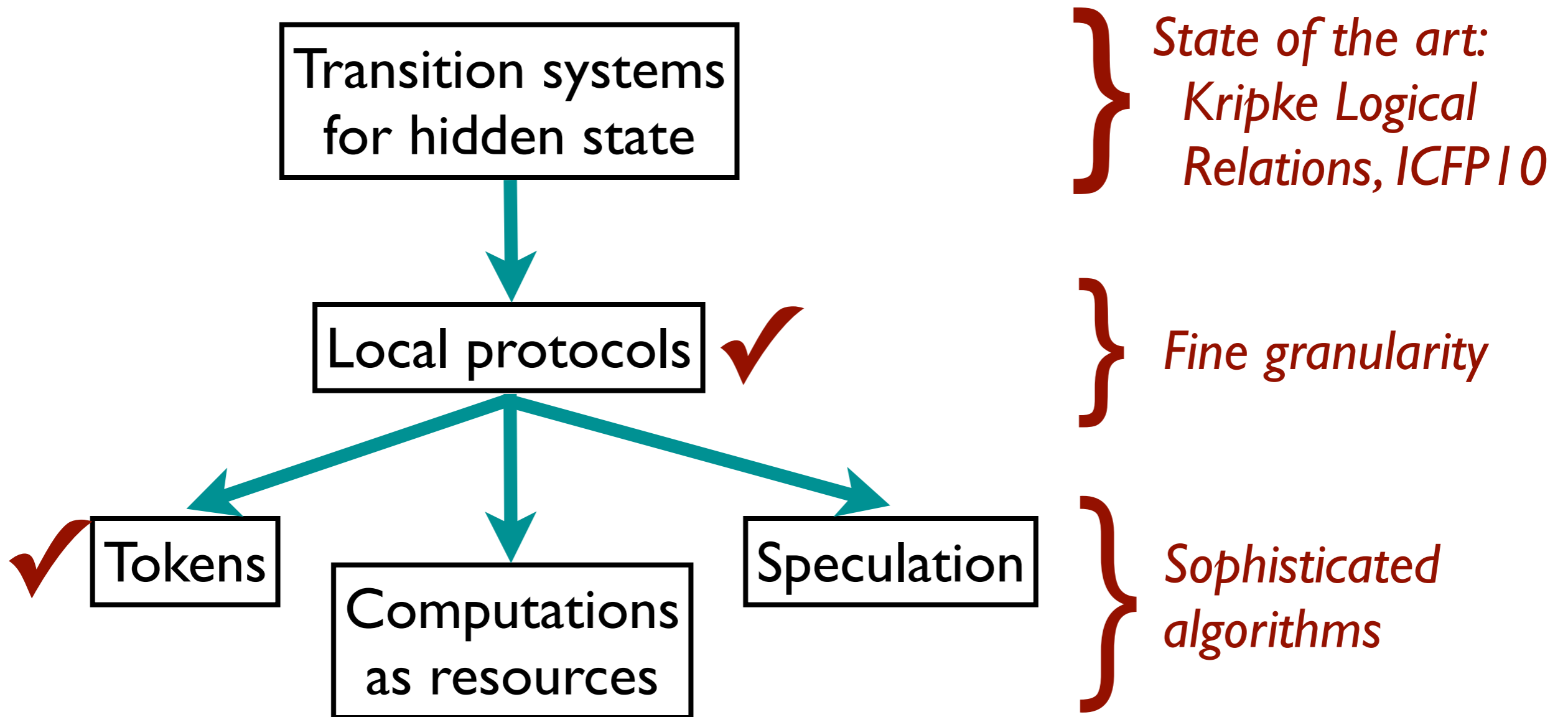
*Rely*: transitions with environment tokens

*Guarantee*: transitions with thread's tokens

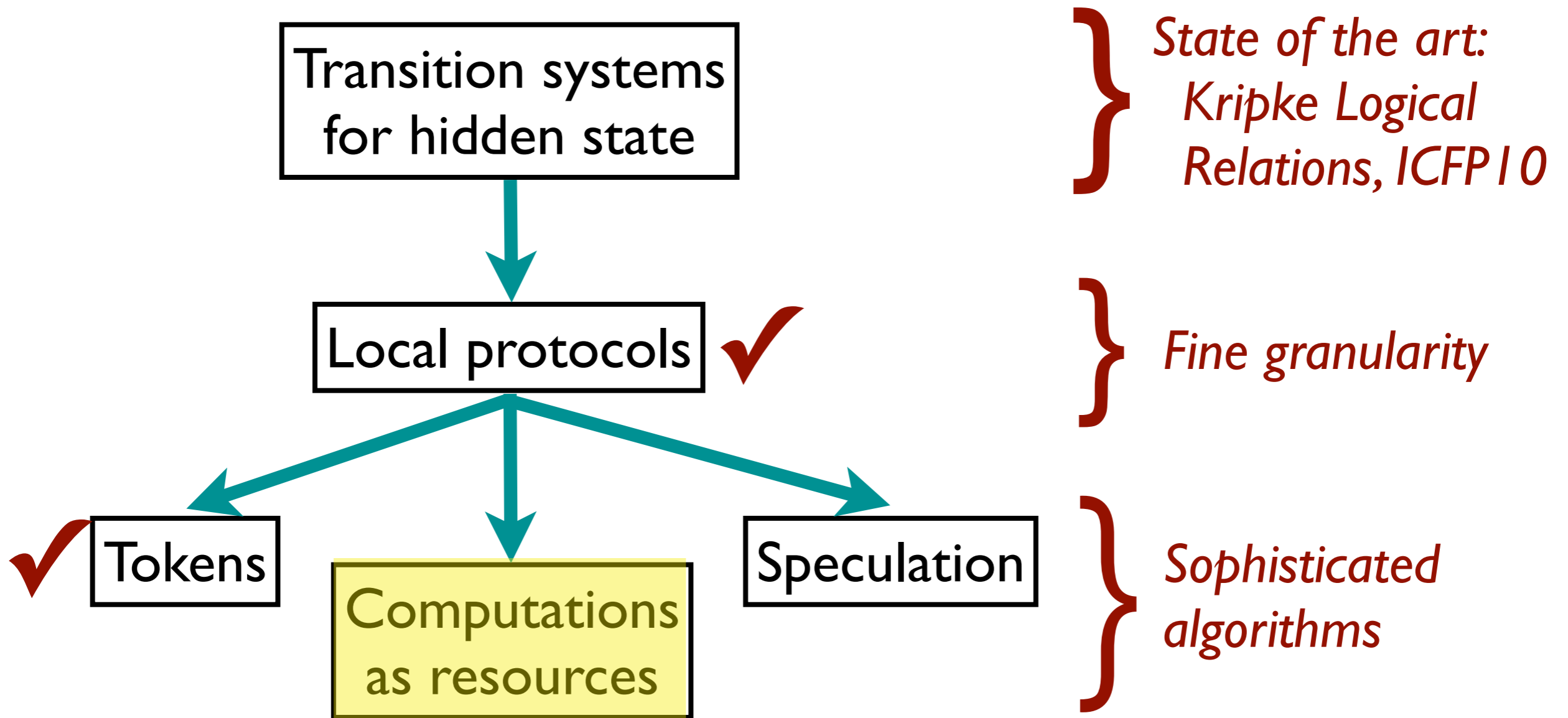
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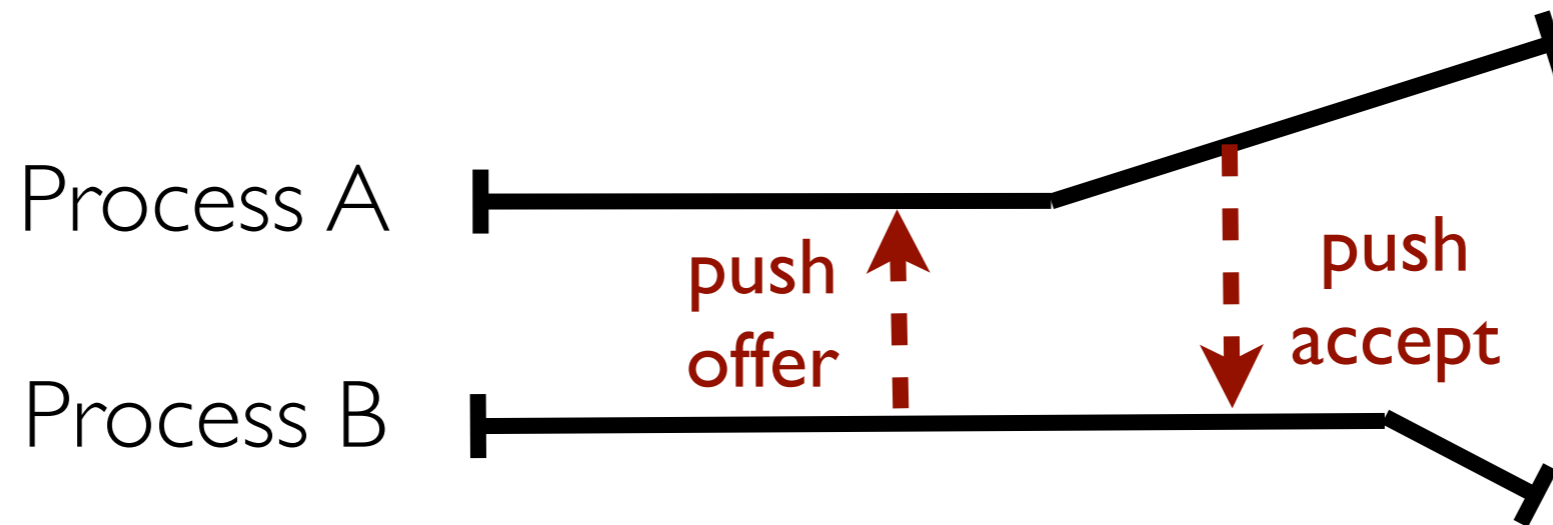
# Approach





## Problem:

Cooperation is inherently non-thread-local  
e.g. elimination stacks

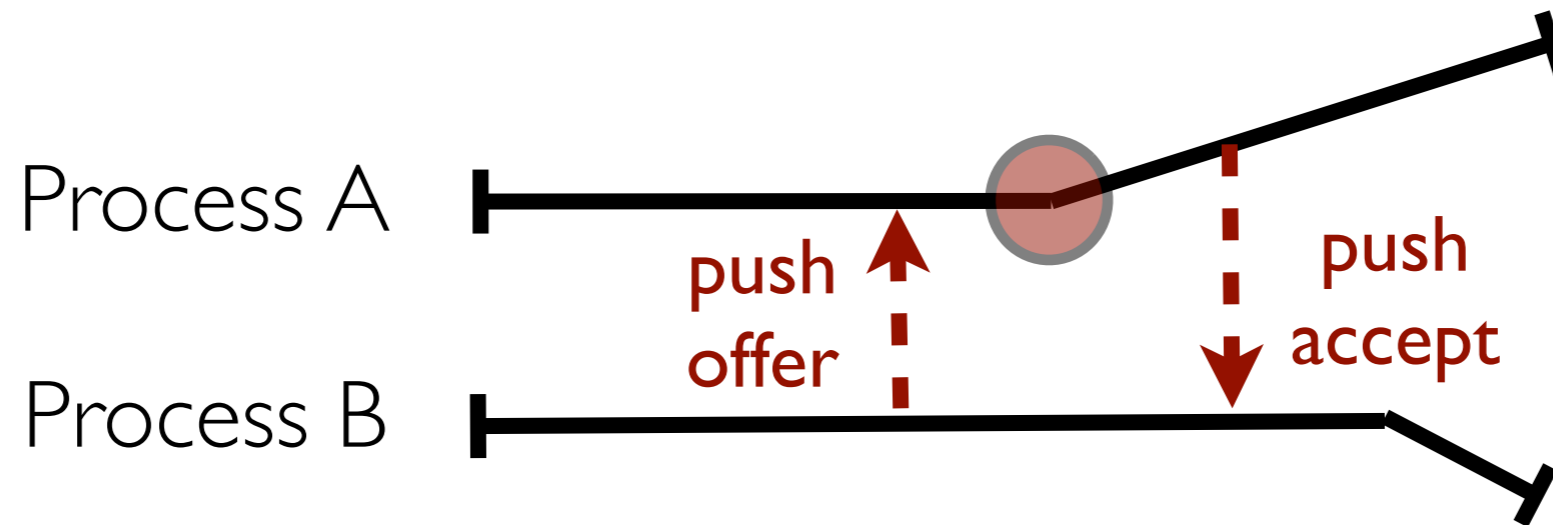


## Solution:

Make thread's specification a *shareable resource*

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## Solution:

Make thread's specification a *shareable resource*

Algorithm	Local protocols	Tokens	Computational resources	Speculation
Treiber's Stack	X			
Michael-Scott Queue	X			
Hand-over-hand set	X	X		
Elimination Flags	X	X	X	
Lazy set	X	X	X	X
Conditional CAS	X	X	X	X
K-CAS	X	X	X	X



# The Takeaway

Analyze fine-grained algorithms  
with fine-grained protocols!

To scale up, must address:

- varying roles
- cooperation
- nondeterminism

Our contributions:

- tokens
- computational resources
- speculation